

# Limitations of quantum walks and search

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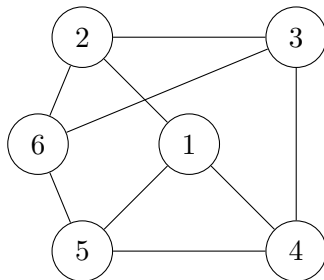
Joint Estonian-Latvian Theory Days

# Outline

- 1 Classical Random Walks
- 2 Quantum Walks
- 3 Localization
- 4 Quantum Search
- 5 Stationary States

# Classical Random Walks

- We walk on an  $N$  vertex graph.



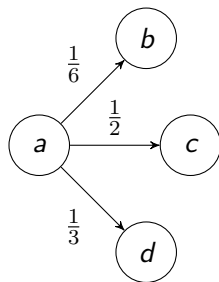
## Classical Random Walks

- A state of the walk  $x$  is a probability distribution over the vertices –  $x(i)$  is the probability of being at vertex  $i$ .

- $$\sum_{i \in [N]} x(i) = 1.$$

- The probability of going from  $u$  to  $v$  in a single step:  $p_{uv}$ .

- $$\sum_{u \sim v} p_{uv} = 1.$$



## Walk Operator

- Let  $x_t$  – probability distribution after  $t$  steps. We start at  $x_0$ .
- A single step of the walk – a matrix operator:

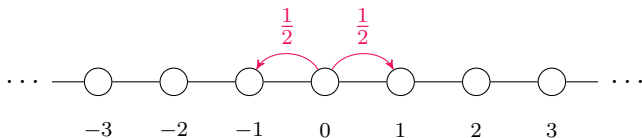
$$P = \begin{bmatrix} p_{11} & p_{21} & p_{31} & \dots & p_{N1} \\ p_{12} & p_{22} & p_{32} & \dots & p_{N2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ p_{1N} & p_{2N} & p_{3N} & \dots & p_{NN} \end{bmatrix}$$

- We can express  $t$  steps as  $P^t$ :

$$x_t = P x_{t-1} = P^t x_0.$$

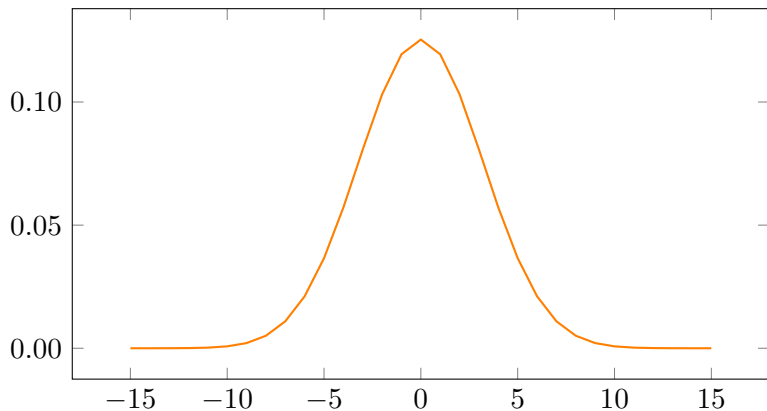
## Example: Walk on a Line

- $x_0(0) = 1$ .
- $p_{i,i-1} = p_{i,i+1} = \frac{1}{2}$ .



## Example: Walk on a Line

- Probability after 40 steps.



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# Discrete Time Quantum Walk

- Instead of vertices, our state is distributed over directed edges.
- For every edge  $u \rightarrow v$  we have a basis state  $|uv\rangle$ .
- A state of the walk – a vector over the edges:

$$|\psi\rangle = \sum_{u \rightarrow v} \alpha_{uv} |uv\rangle.$$

## Discrete Time Quantum Walk

- Instead of probabilities,  $\alpha_{uv}$  are complex amplitudes.
- The probability of obtaining  $|uv\rangle$  is  $|\alpha_{uv}|^2$ .
- $\sum_{u \rightarrow v} |\alpha_{uv}|^2 = 1$ .



## Quantum Walk Operator

- $|\psi_t\rangle$  – the state after  $t$  steps. We start at  $|\psi_0\rangle$ .
- A single step of the walk – some unitary transformation  $U$ .

$$|\psi_t\rangle = U|\psi_{t-1}\rangle = U^t|\psi_0\rangle.$$

- Afterwards, we measure the state of the walk: with probability  $|\alpha_{uv}(\psi_t)|^2$  we obtain  $u \rightarrow v$ .
- With probability  $\sum_{v:u \rightarrow v} |\alpha_{uv}(\psi_t)|^2$  we obtain position  $u$ .

# Quantum Walk Operator

- $U = SC$
- $C$  – the “coin” operator. This disperses the amplitudes among the directions within a single vertex.
- $S$  – the shift operator. This moves the amplitudes along the edges of the graph.

## Example: Quantum Walk on a Line

- $U = SC$

- $C|x, \leftarrow\rangle = \frac{1}{\sqrt{2}}|x, \leftarrow\rangle + \frac{i}{\sqrt{2}}|x, \rightarrow\rangle$

- $C|x, \rightarrow\rangle = \frac{i}{\sqrt{2}}|x, \leftarrow\rangle + \frac{1}{\sqrt{2}}|x, \rightarrow\rangle$

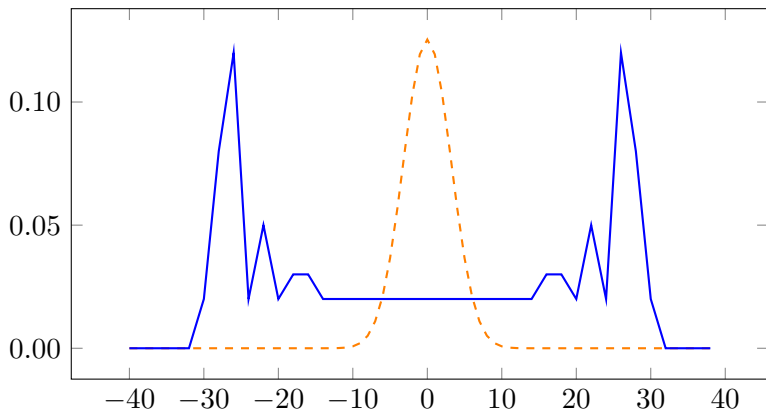
- $S|x, \leftarrow\rangle = |x-1, \leftarrow\rangle$

- $S|x, \rightarrow\rangle = |x+1, \rightarrow\rangle$

- Start at  $|\psi_0\rangle = \frac{1}{\sqrt{2}}|0, \leftarrow\rangle + \frac{1}{\sqrt{2}}|0, \rightarrow\rangle$ .

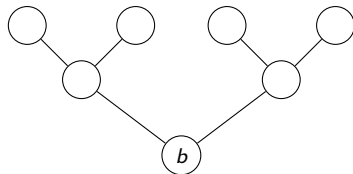
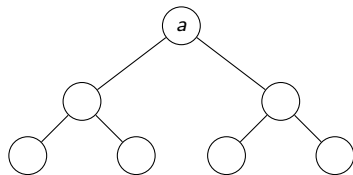
## Example: Quantum Walk on a Line

- Probability after 40 steps.



## Exponential Speedup

- Join two binary trees of depth  $D$  at the leaves.
- The number of steps needed to go from  $a$  to  $b$  with some constant probability:
  - Random walk:  $\mathcal{O}(2^D)$ .
  - Quantum walk:  $\mathcal{O}(D^2)$ .





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# Localization

- *Localization* – a phenomenon where the walk remains at the starting position  $|\psi_0\rangle$  with high probability.
- Formally, for any  $t$ :

$$|\langle\psi_0|U^t|\psi_0\rangle|^2 \approx 1,$$

where  $|\langle a|b\rangle|^2$  is essentially the inner product of  $a$  and  $b$ .

## The Examined Walks

- We examine walks on  $d$ -regular  $N$  vertex graphs.
- As  $C$  we use Grover's diffusion:

$$C|uv\rangle = -|uv\rangle + \frac{2}{d} \sum_{u \rightarrow w} |uw\rangle.$$

An inversion about the average amplitude at each vertex.

- As  $S$  we use the “flip-flop” shift:

$$S|uv\rangle = |vu\rangle.$$

## Our Paper

- A. Ambainis, K. Prūsis, J. Vihrovs, and T. G. Wong.  
Oscillatory localization of quantum walks by classical electric circuits, 2016
- Oscillatory localization – the walk jumps back and forth between two states.
- Formally, for any  $t$ :

$$|\langle \psi_0 | \mathcal{U}^{2t} | \psi_0 \rangle|^2 \approx 1.$$

## Oscillatory Localization Example

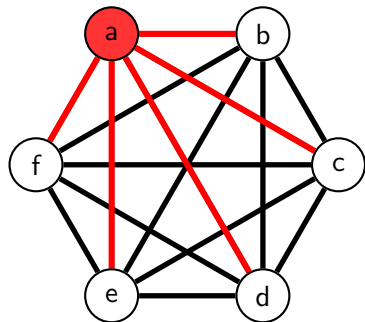
- The complete graph  $K_N$ .

- Starting in the state

$$|\psi_0\rangle = \frac{1}{\sqrt{N-1}} \sum_{a \rightarrow v} |av\rangle,$$

the walk disperses.

- Starting in  $|\psi_0\rangle = |ab\rangle$ , it localizes.



## Oscillatory Localization Example

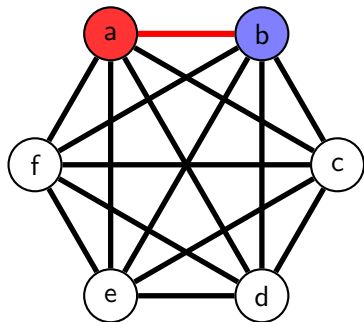
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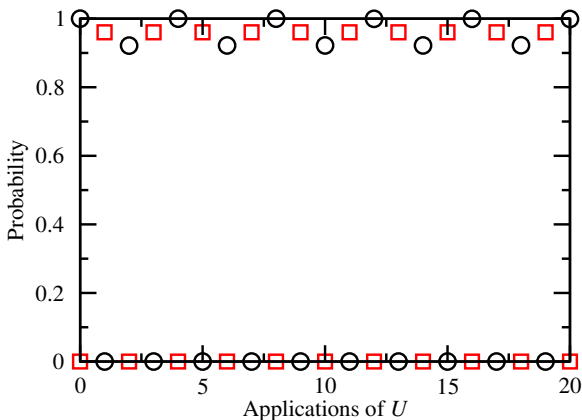
the walk disperses.

- Starting in  $|\psi_0\rangle = |ab\rangle$ , it localizes.



## Oscillatory Localization Example

- Example for  $N = 16$ , black circles are probability at  $|ab\rangle$ , red squares are probability at  $|ba\rangle$ .



# Eigenvectors

- In fact, we can look at the 1-eigenvectors of  $U^2$ .
- The walk localizes if  $|\psi_0\rangle$  is close to these.
- It turns out there are only two types of such eigenvectors.



## Uniform States

- The first type we call *uniform states*.
- If the graph is not bipartite, the only such state is the uniform distribution over all edges:

$$|\sigma\rangle = \sum_{u \rightarrow v} \frac{1}{\sqrt{dN}} |uv\rangle.$$

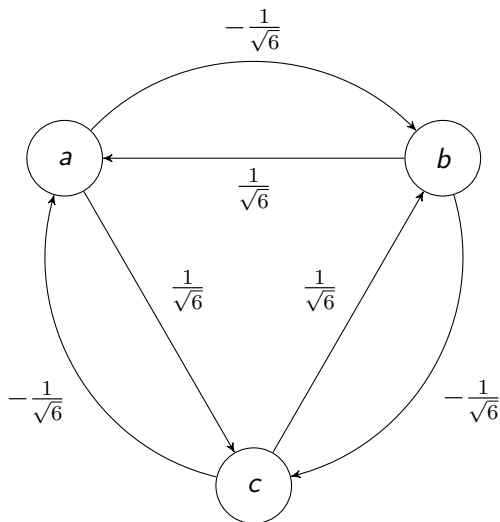
- If the graph is bipartite, these are uniform over the edges of each part.

## Flip States

- We call the second type *flip states*.
- A state  $|\phi\rangle$  is a flip state if the following two conditions hold for every vertex  $v$ :

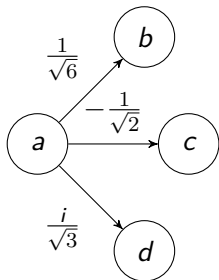
$$\sum_{u \rightarrow v} \alpha_{uv} = 0, \quad \sum_{u \leftarrow v} \alpha_{vu} = 0.$$

# Flip State Example

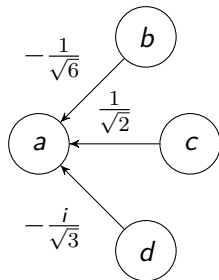


## Flipped States

■  $|\psi\rangle$



■  $|\tilde{\psi}\rangle$



■ For flip states,  $U|\phi\rangle = |\tilde{\phi}\rangle$ .

## Our Result

- Express the starting state in terms of normalized flip and uniform states with some remainder state  $|\rho\rangle$ :

$$|\psi_0\rangle = \alpha |\phi\rangle + \beta |\sigma\rangle + \gamma |\rho\rangle.$$

- After an even number of steps  $2t$ ,

$$|\langle\psi_0|U^{2t}|\psi_0\rangle| \geq 2(|\alpha|^2 + |\beta|^2) - 1.$$

- After an odd number of steps  $2t + 1$ ,

$$|\langle\widetilde{\psi}_0|U^{2t+1}|\psi_0\rangle| \geq 2 \max(|\alpha|^2, |\beta|^2) - 1.$$

## Projection unto Flip States

- We want to find flip states close to our starting state.
- We do not know how to do this for arbitrary starting states.
- For single edge starting states, we can do this using electric networks.

## Result

- High connectivity between neighboring vertices  $a, b$  implies oscillatory localization single edge starting states:
  - $|ab\rangle$ ,
  - $\frac{1}{\sqrt{2}}|ab\rangle - \frac{1}{\sqrt{2}}|ba\rangle$ .
- Many common graphs have high connectivity, such as high-degree edge transitive graphs, which include the complete graph and the hypercube.

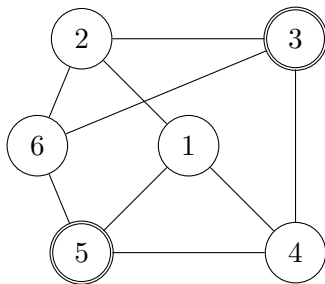
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## Search on a Graph

- We walk on an  $N$  vertex graph, where some vertices are marked.



## Quantum Search

- Denote the number of edges of the graph by  $|E|$ .
- We start in a uniform superposition over all edges of the graph:

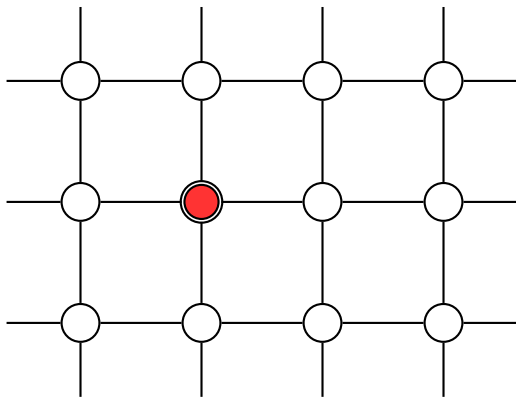
$$|\psi_0\rangle = \frac{1}{\sqrt{2|E|}} \sum_{u \rightarrow v} |uv\rangle.$$

- $U = SCQ$ .

$$Q|uv\rangle = \begin{cases} |uv\rangle & u \text{ is not marked;} \\ -|uv\rangle & u \text{ is marked.} \end{cases}$$

## Example: Quantum Search on a Grid

- This algorithm finds a marked vertex on the periodic 2D lattice in  $\mathcal{O}(\sqrt{N} \log N)$  time.



# Outline

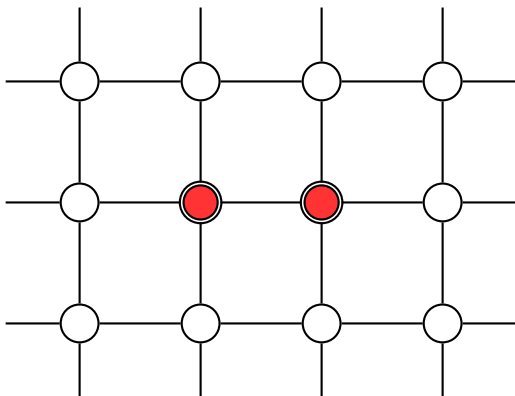
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## Multiple Marked Vertices

- Classically, additional marked vertices always make search easier – the expected runtime for  $k$  marked vertices is  $\mathcal{O}(N/k)$ .
- This is not always the case for quantum search.
- In some cases, multiple marked vertices can make the search remain at the starting uniform state.

## Multiple Marked Vertices Example

- Nahimov's result: quantum search does not work on the 2D lattice with 2 adjacent marked vertices.



## Stationary States

- In this case, we look at the 1-eigenvectors of  $U = SCQ$ .
- The search remains stationary if the uniform starting state

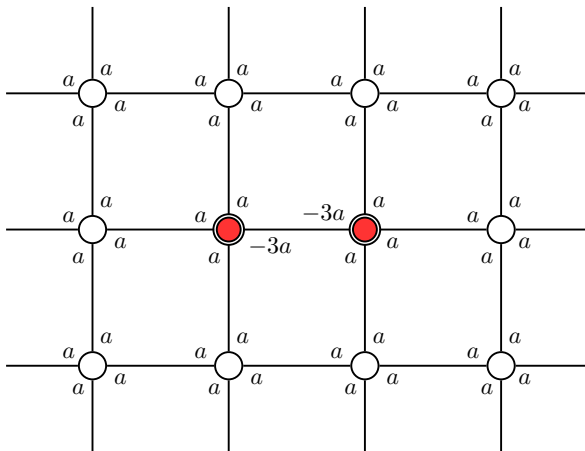
$$|\psi_0\rangle = \frac{1}{\sqrt{2|E|}} \sum_{u \rightarrow v} |uv\rangle.$$

is close to some stationary  $|\psi\rangle$ :

$$|\langle\psi_0|\psi\rangle|^2 \approx 1.$$

## Stationary State Example

- Nahimov's stationary state:





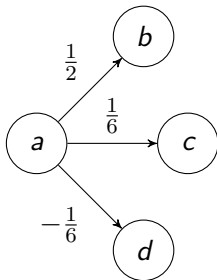
## Our Paper

- K. Prūsis, J. Vihrovs, and T. G. Wong. [Stationary states in quantum walk search](#).  
*Phys. Rev. A*, 94:032334, Sep 2016
- We give some general criteria for when such stationary states exist.
- We give additional criteria for stationary states close to the starting uniform states.

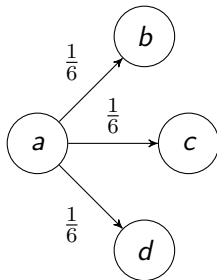
## Flip and Uniform Components

- For any state  $|\psi\rangle$ , we can decompose its amplitudes at a vertex  $|\psi_v\rangle$  into uniform and flip components:

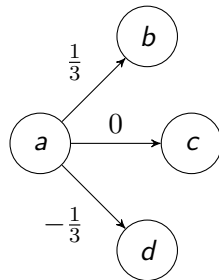
- $|\psi_a\rangle$



- $|\sigma_a\rangle$



- $|\phi_a\rangle$



## Search Operator on Flip and Uniform Components

- If  $v$  is not marked:

$$CQ|\sigma_v\rangle = |\sigma_v\rangle \quad CQ|\phi_v\rangle = -|\phi_v\rangle.$$

- If  $v$  is marked:

$$CQ|\sigma_v\rangle = -|\sigma_v\rangle \quad CQ|\phi_v\rangle = |\phi_v\rangle.$$

## Stationary State Conditions

- A state is stationary if and only if, for every edge  $ab$ :
  - If exactly one of  $a, b$  is marked, the following holds (for  $b$  marked):

$$\sigma_{ab} = \phi_{ba} \quad \phi_{ab} = -\sigma_{ba}.$$

- If they are both unmarked:

$$\sigma_{ab} = \sigma_{ba} \quad \phi_{ab} = -\phi_{ba}.$$

- If they are both marked:

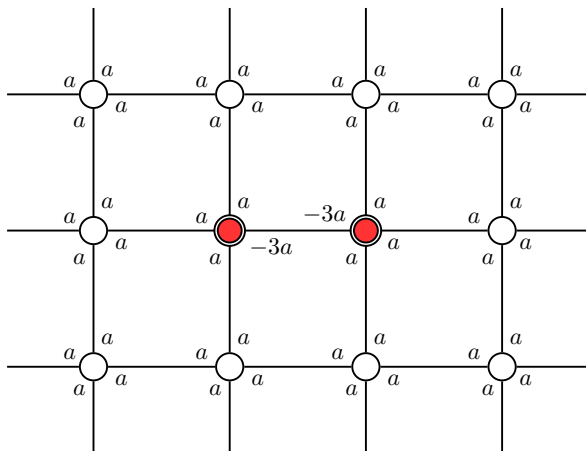
$$\sigma_{ab} = -\sigma_{ba} \quad \phi_{ab} = \phi_{ba}.$$

## Optimal Stationary States

- We are interested in finding the stationary state closest to the starting state.
- The flip components of unmarked vertices and the uniform components of marked vertices contribute 0 to  $|\langle \psi_0 | \psi \rangle|^2$ .
- Theorem: the stationary state  $|\psi\rangle$  maximizing  $|\langle \psi_0 | \psi \rangle|^2$  satisfies the following conditions:
  - If  $v$  is unmarked,  $|\psi_v\rangle = |\sigma_v\rangle$ .
  - If  $v$  is marked,  $|\psi_v\rangle = |\phi_v\rangle$ .
  - For every adjacent  $u, v$ ,  $\alpha_{uv} = \alpha_{vu}$ .

## Existing Examples are Optimal

- This matches the conditions for the grid.



## Existence Criteria

- A stationary state exists if every connected component of marked vertices is either:
  - *not* bipartite;
  - bipartite, but the sums of the amplitudes on the edges going to unmarked vertices are equal for each part.

# Questions?