

Tutorial on Quantum Computing

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Eiropas Sociālā fonda projekts
“Datorzinātnes pielietojumi un tās saiknes ar kvantu fiziku”
Nr.2009/0216/1DP/1.1.1.2.0/09/APIA/VIAA/044

Outline of the tutorial

- Prehistory and history of quantum computing
- Memory
 - Qubit
 - Qubit register
- Computation: unitary operators
- Reading outcome: measurement
- Efficient quantum algorithms
 - Deutsch-Jozsa algorithm
 - Grover's algorithm



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Prehistory: quantum mechanics

- 1900: quantum hypothesis of Max Planck
- 1905: Albert Einstein's postulation of light quanta (photons)
- 1913: Niels Bohr's model of atomic structure
- 1924: Louis de Broglie's hypothesis of wave-particle duality
- 1925: Wolfgang Pauli's exclusion principle
- 1926: Erwin Schrödinger's equation
- 1926: probability density function by Max Born
- 1927: Werner Heisenberg's uncertainty principle
- 1928: Paul Dirac's equation
- 1932: mathematical foundations of QM by John von Neumann
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Prehistory: theory of computation

- 1931: Kurt Gödel's incompleteness theorems
- 1936: Emil Leon Post's machine
- 1936: undecidable problem, λ -calculus by Alonzo Church
- 1936: undecidable problem, Turing machine by Alan Turing
- 1939: Stephen Cole Kleene's recursion theory
- 1945: John von Neumann's computer architecture
- 1948: information theory by Claude E. Shannon
- 1956: Noam Chomsky's grammar hierarchy
- 1965: complexity theory, Juris Hartmanis and Richard Stearns
- 1971: Stephen Cook formulates the $P = NP?$ problem
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History of quantum computation

- 1973: reversible computation, Charles H. Bennett
- 1973: Alexander Holevo's bound on quantum information
- 1981: Richard Feynman's idea of a quantum computer
- 1984: quantum key distribution, Charles H. Bennett and Gilles Brassard
- 1985: universal quantum computer by David Deutsch
- 1993: Dan Simon's algorithm: exponential speed-up in an oracle problem
- 1994: Peter Shor's quantum poly-time factoring algorithm
- 1996: Lov Grover's quadratic speed-up database search
- 1998: first small quantum computers
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Basic idea

- Computation must be performed as a real, physical process, therefore it must obey physical laws
- At the fundamental level the physics are described by quantum mechanics
- Does quantum mechanics imply any differences to the (classical) computation as we know it? – YES!



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Memory: a classical bit

- Is either 0 or 1
- A one-bit-memory in the classical sense is a (classical-physics) system with two possible distinguishable states designated by 0 and 1



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Memory: a quantum bit (qubit)

- A two-state quantum system
- Its two distinguishable states are designated by $|0\rangle$ and $|1\rangle$



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Quantum principle No. 1

- Quantum principle of superposition: if a quantum system can be in any of n distinguishable (basis) states, then it can also be in any **superposition** of these states, with complex numbers $\alpha_1, \alpha_2, \dots, \alpha_n$, called (probability) **amplitudes**, characterizing the amount by which the system is in respective basis states. It must hold:

$$|\alpha_1|^2 + |\alpha_2|^2 + \dots + |\alpha_n|^2 = 1$$



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Memory: a quantum bit (qubit)

- Thus qubit can be also in any superposition of $|0\rangle$ and $|1\rangle$, with some amplitudes α_0 and α_1
- Such superposition state is denoted:

$$\alpha_0 |0\rangle + \alpha_1 |1\rangle$$

- α_0 and α_1 are complex numbers with

$$|\alpha_0|^2 + |\alpha_1|^2 = 1$$



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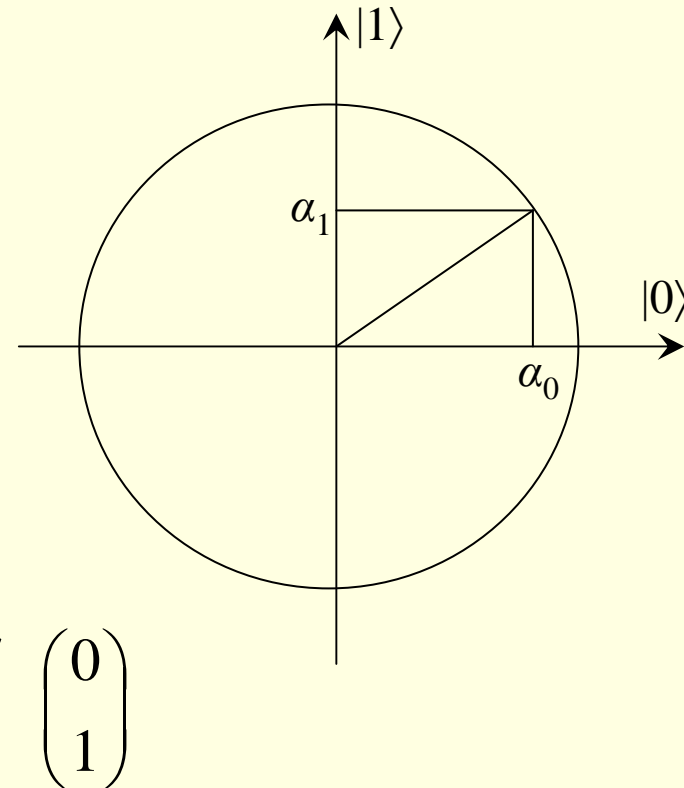
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Memory: a quantum bit (qubit)

- The state of a qubit can also be described by a vector in \mathbb{C}^2 : $\begin{pmatrix} \alpha_0 \\ \alpha_1 \end{pmatrix}$

- Thus $|0\rangle = 1|0\rangle + 0|1\rangle$ or $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$

- ... and $|1\rangle = 0|0\rangle + 1|1\rangle$ or $\begin{pmatrix} 0 \\ 1 \end{pmatrix}$



Memory: a quantum bit (qubit)

- The exact state of an unknown qubit $\alpha_0 |0\rangle + \alpha_1 |1\rangle$ **cannot be learned**
- The information from a qubit can be obtained by a **measurement** (we will talk about it in more detail later)
- One can measure whether qubit is in state $|0\rangle$ or $|1\rangle$
- Then one will obtain:
 - the answer $|0\rangle$ with probability $|\alpha_0|^2$,
 - the answer $|1\rangle$ with probability $|\alpha_1|^2$
- After the measurement the qubit collapses to the state equal with the given answer

Memory: a qubit register

- We can put 2 qubits $\alpha_0 |0\rangle + \alpha_1 |1\rangle$ and $\beta_0 |0\rangle + \beta_1 |1\rangle$ together forming a two qubit register
- Then if we measure both qubits, we will obtain for the first qubit:
 - the answer $|0\rangle$ with probability $|\alpha_0|^2$,
 - the answer $|1\rangle$ with probability $|\alpha_1|^2$
- ... and for the second qubit:
 - the answer $|0\rangle$ with probability $|\beta_0|^2$,
 - the answer $|1\rangle$ with probability $|\beta_1|^2$

Memory: a qubit register

- Alternatively, we can look on a two qubit register as on a four state $|0\rangle|0\rangle$, $|0\rangle|1\rangle$, $|1\rangle|0\rangle$, $|1\rangle|1\rangle$ system
- From such viewpoint we will get:
 - the answer $|0\rangle|0\rangle$ with probability $|\alpha_0\beta_0|^2$,
 - the answer $|0\rangle|1\rangle$ with probability $|\alpha_0\beta_1|^2$,
 - the answer $|1\rangle|0\rangle$ with probability $|\alpha_1\beta_0|^2$,
 - the answer $|1\rangle|1\rangle$ with probability $|\alpha_1\beta_1|^2$

Memory: a qubit register

- A multiplication law works:

$$\begin{aligned} & (\alpha_0|0\rangle + \alpha_1|1\rangle) \otimes (\beta_0|0\rangle + \beta_1|1\rangle) = \\ & \alpha_0\beta_0|0\rangle \otimes |0\rangle + \alpha_0\beta_1|0\rangle \otimes |1\rangle + \alpha_1\beta_0|1\rangle \otimes |0\rangle + \alpha_1\beta_1|1\rangle \otimes |1\rangle \end{aligned}$$

- $|0\rangle \otimes |0\rangle$ denotes the same as $|0\rangle |0\rangle$ or $|00\rangle$
- This multiplication is called **tensor multiplication**
- It is not commutative:

$|0\rangle \otimes |1\rangle$ is not the same as $|1\rangle \otimes |0\rangle$

Memory: a qubit register

- A **tensor** or **Kronecker product** of matrices:

$A = (a_{ij})$ of size $m \times n$, $B = (b_{ij})$ of size $k \times l$

$$A \otimes B = \begin{pmatrix} a_{11}B & a_{12}B & \dots & a_{1n}B \\ a_{21}B & a_{22}B & \dots & a_{2n}B \\ \dots & \dots & \dots & \dots \\ a_{m1}B & a_{m2}B & \dots & a_{mn}B \end{pmatrix} \text{ of size } mk \times nl$$



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Memory: a qubit register

- For our two qubits in matrix form:

$$\begin{pmatrix} \alpha_0 \\ \alpha_1 \end{pmatrix} \otimes \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix} = \begin{pmatrix} \alpha_0 \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix} \\ \alpha_1 \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} \alpha_0 \beta_0 \\ \alpha_0 \beta_1 \\ \alpha_1 \beta_0 \\ \alpha_1 \beta_1 \end{pmatrix}$$

Quantum principle No. 1

- Quantum principle of superposition: if a quantum system can be in any of n distinguishable (basis) states, then it can also be in any **superposition** of these states, with complex numbers $\alpha_1, \alpha_2, \dots, \alpha_n$, called (probability) **amplitudes**, characterizing the amount by which the system is in respective basis states. It must hold:

$$|\alpha_1|^2 + |\alpha_2|^2 + \dots + |\alpha_n|^2 = 1$$



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Memory: a qubit register

- General state of a two qubit register:

$$\alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$

$$\text{where } |\alpha_{00}|^2 + |\alpha_{01}|^2 + |\alpha_{10}|^2 + |\alpha_{11}|^2 = 1$$



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Memory: a qubit register

- There are such states that cannot be expressed as a tensor product of two 1-qubit states
- It means that the 2-qubit register cannot be looked at as a composition of two independent qubits
- Such states are called **entangled** states



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Memory: a qubit register

- Example, the Bell state: $|\Phi^+\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$
- The 1st qubit:
 - is $|0\rangle$ with prob. $1/2$,
 - is $|1\rangle$ with prob. $1/2$
- The 2nd qubit:
 - is $|0\rangle$ with prob. $1/2$,
 - is $|1\rangle$ with prob. $1/2$
- ... but the register:
 - is $|00\rangle$ with prob. $1/2$,
 - is $|01\rangle$ with prob. 0 ,
 - is $|10\rangle$ with prob. 0 ,
 - is $|11\rangle$ with prob. $1/2$
- The multiplication law does not work!



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Memory: a qubit register

- General state of an n qubit register:

$$\alpha_{0\dots 00}|0\dots 00\rangle + \alpha_{0\dots 01}|0\dots 01\rangle + \dots + \alpha_{1\dots 11}|1\dots 11\rangle$$

$$\text{where } |\alpha_{0\dots 00}|^2 + |\alpha_{0\dots 01}|^2 + \dots + |\alpha_{1\dots 11}|^2 = 1$$

- Geometrically: an arbitrary vector of unit length in \mathbb{C}^{2^n}

$$\begin{pmatrix} \alpha_{0\dots 00} \\ \alpha_{0\dots 01} \\ \dots \\ \alpha_{1\dots 11} \end{pmatrix}$$

Quantum principle No. 2

- Quantum principle of state evolution: the change of the state of a quantum system is a **unitary** linear operator
- A linear operator is unitary iff it preserves the vector norm
- ... alternatively, it maps the unit hypersphere (where the quantum state vectors reside) to itself
- Essentially, unitary operators are the rotations
- Since unitary operators are linear, to define them it is enough to specify their action on the basis vectors



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Computation: unitary operators

- 1-qubit unitary operator examples:

identity operator : $id = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$

NOT operator : $NOT = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

rotation : $R_\varphi = \begin{pmatrix} \cos \varphi & -\sin \varphi \\ \sin \varphi & \cos \varphi \end{pmatrix}$

phase shift : $P_\varphi = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\varphi} \end{pmatrix}$

Computation: unitary operators

- 1-qubit unitary operator example, the Hadamard transform:

$$|0\rangle \rightarrow \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$

$$|1\rangle \rightarrow \frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle$$

$$H = \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}$$

Computation: unitary operators

- 2-qubit unitary operator example, the controlled NOT operator:

$$|00\rangle \rightarrow |00\rangle$$

$$|01\rangle \rightarrow |01\rangle$$

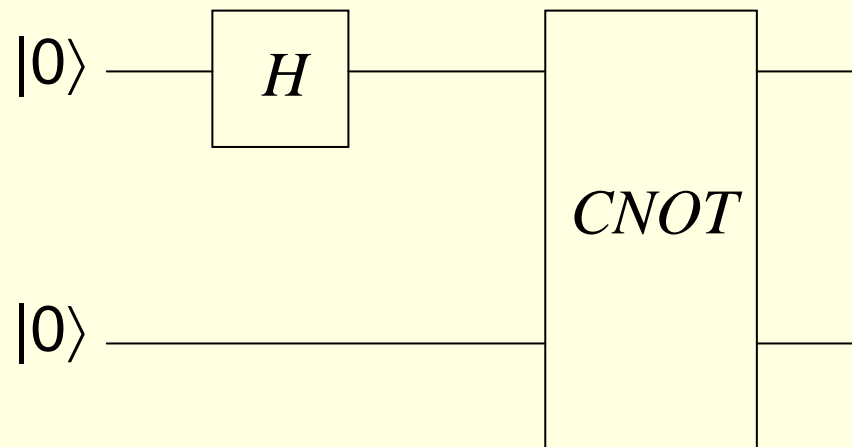
$$|10\rangle \rightarrow |11\rangle$$

$$|11\rangle \rightarrow |10\rangle$$

$$CNOT = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

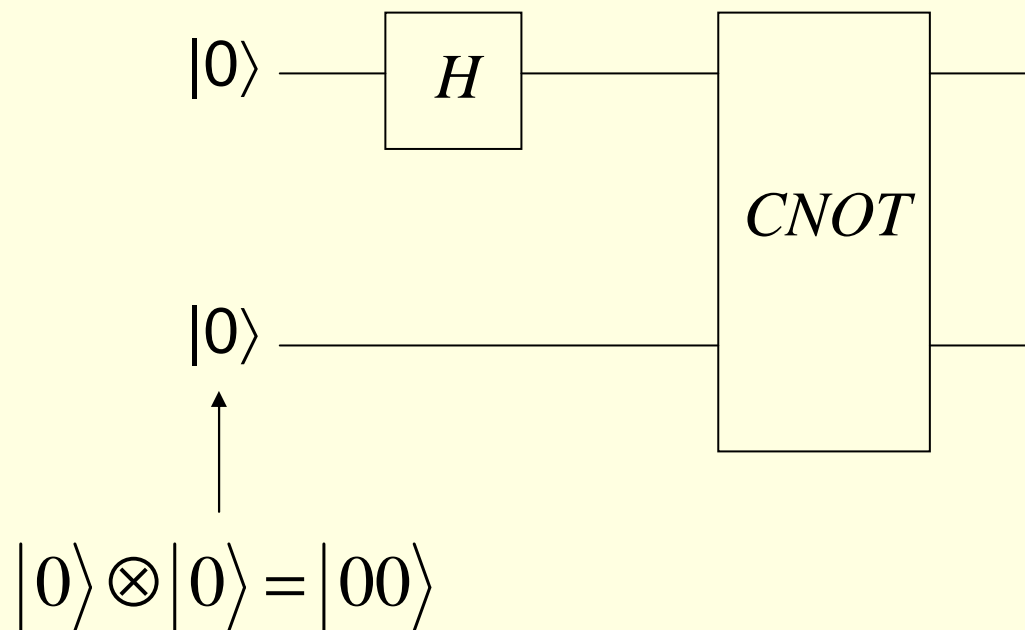
Computation: unitary operators

- Example of a computation, creation of the Bell state $|\Phi^+\rangle$:



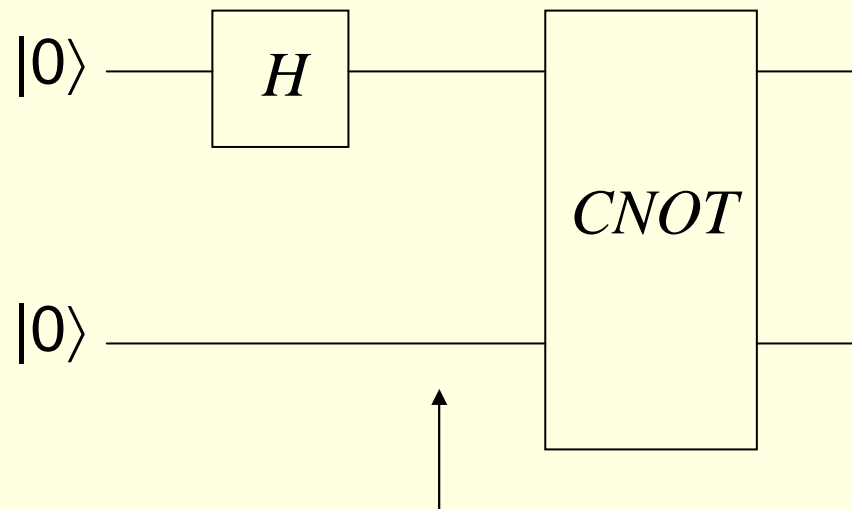
Computation: unitary operators

- Example of a computation, creation of the Bell state $|\Phi^+\rangle$:



Computation: unitary operators

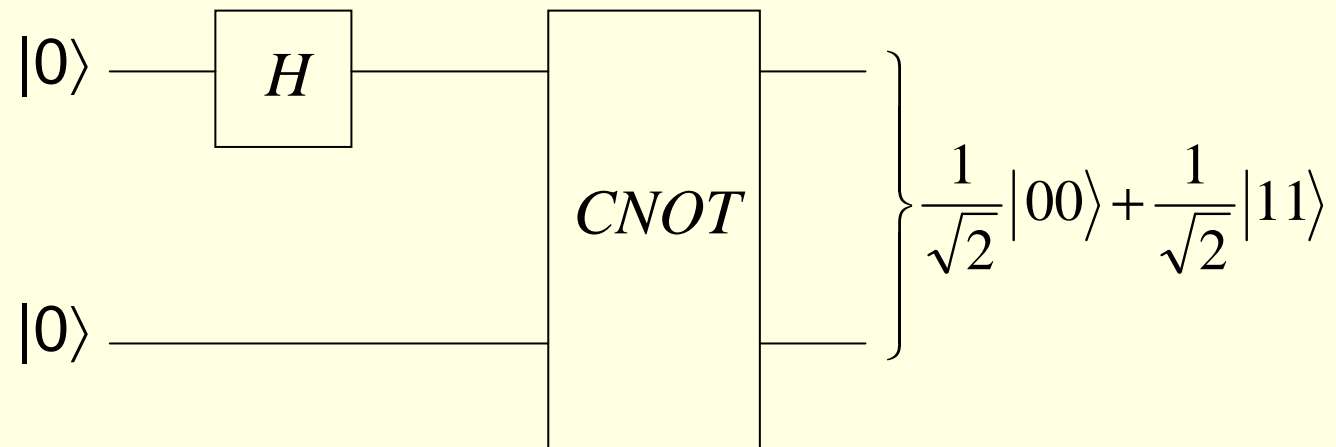
- Example of a computation, creation of the Bell state $|\Phi^+\rangle$:



$$\left(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle \right) \otimes |0\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|10\rangle$$

Computation: unitary operators

- Example of a computation, creation of the Bell state $|\Phi^+\rangle$:



Quantum principle No. 3

- Any information about the state of a quantum system can be extracted into the macroscopic world only by means of a **measurement**
- Mathematically measurement means partitioning the state space H into orthogonal subspaces: $H = E_1 \oplus E_2 \oplus \dots \oplus E_k$
- If the state vector before measurement is $|\psi\rangle = \sum_{i=1}^k |proj_{E_i} \psi\rangle$
- ... then after the measurement the state collapses randomly, with probability $\|proj_{E_i} \psi\|^2$ to one of the subspaces:

$$|\psi\rangle \rightarrow |proj_{E_i} \psi\rangle / \|proj_{E_i} \psi\|$$

- The only classical information obtained is i



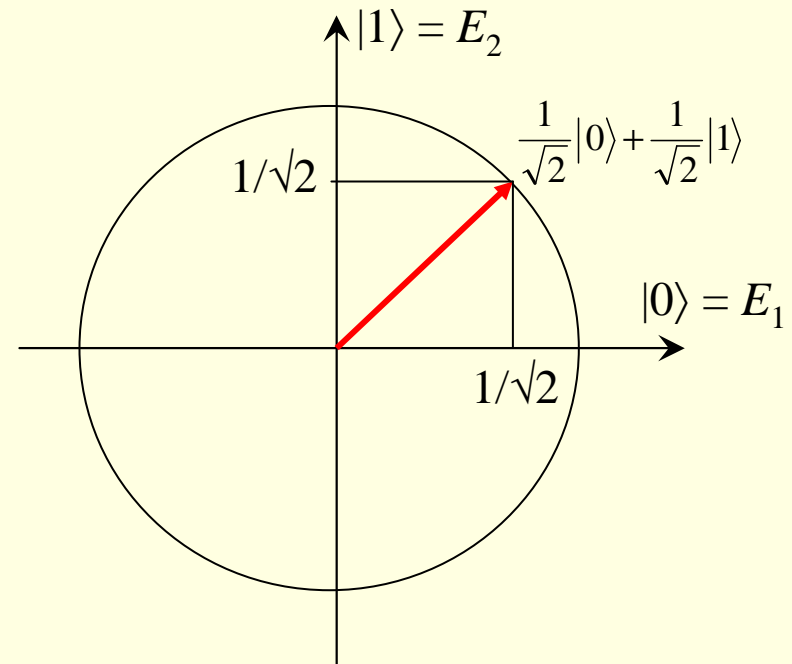
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Reading outcome: measurement

- Example of a measurement:

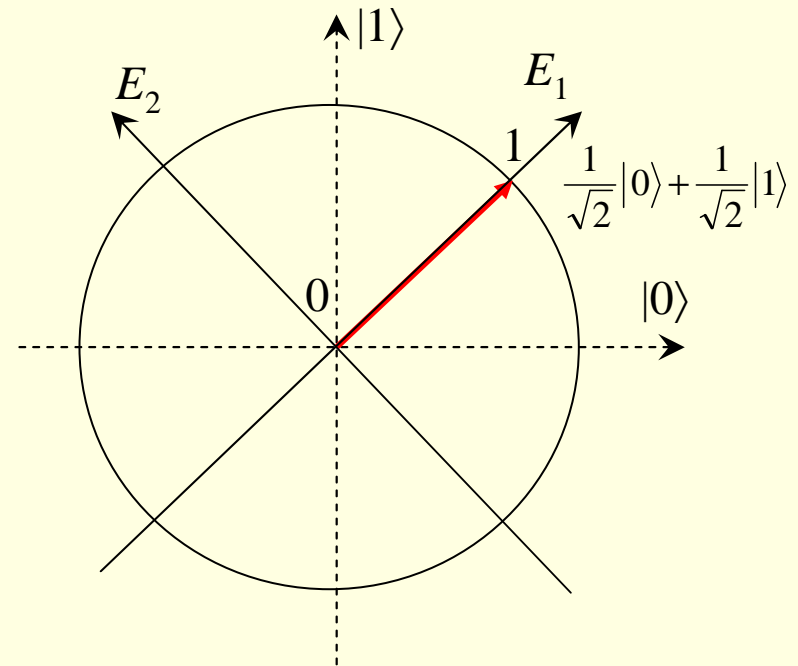


- Outcome E_1 with probability $1/2$
- Outcome E_2 with probability $1/2$

Reading outcome: measurement

- Example of a measurement:

- Outcome E_1 with probability 1
- Outcome E_2 with probability 0



Reading outcome: measurement

- Example: measuring only the first qubit of a 2-qubit system

$$|\psi\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$$

$$E_1 = \text{span}(|00\rangle, |01\rangle)$$

$$E_2 = \text{span}(|10\rangle, |11\rangle)$$

Outcome E_1 with probability $\frac{1}{2}$

Outcome E_2 with probability $\frac{1}{2}$



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Reading outcome: measurement

- Example: measuring only the first qubit of a 2-qubit system

$$|\psi\rangle = \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$$

$$E_1 = \text{span}\left(\left(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle\right) \otimes |0\rangle, \left(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle\right) \otimes |1\rangle\right)$$

$$E_2 = \text{span}\left(\left(\frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle\right) \otimes |0\rangle, \left(\frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle\right) \otimes |1\rangle\right)$$

Outcome E_1 with probability $\frac{1}{2}$

Outcome E_2 with probability $\frac{1}{2}$



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Efficient quantum algorithm

- Example: the Deutsch-Jozsa algorithm (1992)
- Input: a black-box function $f: \{0,1\}^n \rightarrow \{0,1\}$ which is
 - either constant (i.e. all its values are equal),
 - or balanced (i.e. half of its values are 0, and the other half are 1)
- The algorithm can query the black box; the number of queries determines the complexity of algorithm
- The black box works like this: input $|x\rangle|b\rangle$, output $|x\rangle|b \oplus f(x)\rangle$
- Output: answer “constant” or “balanced”



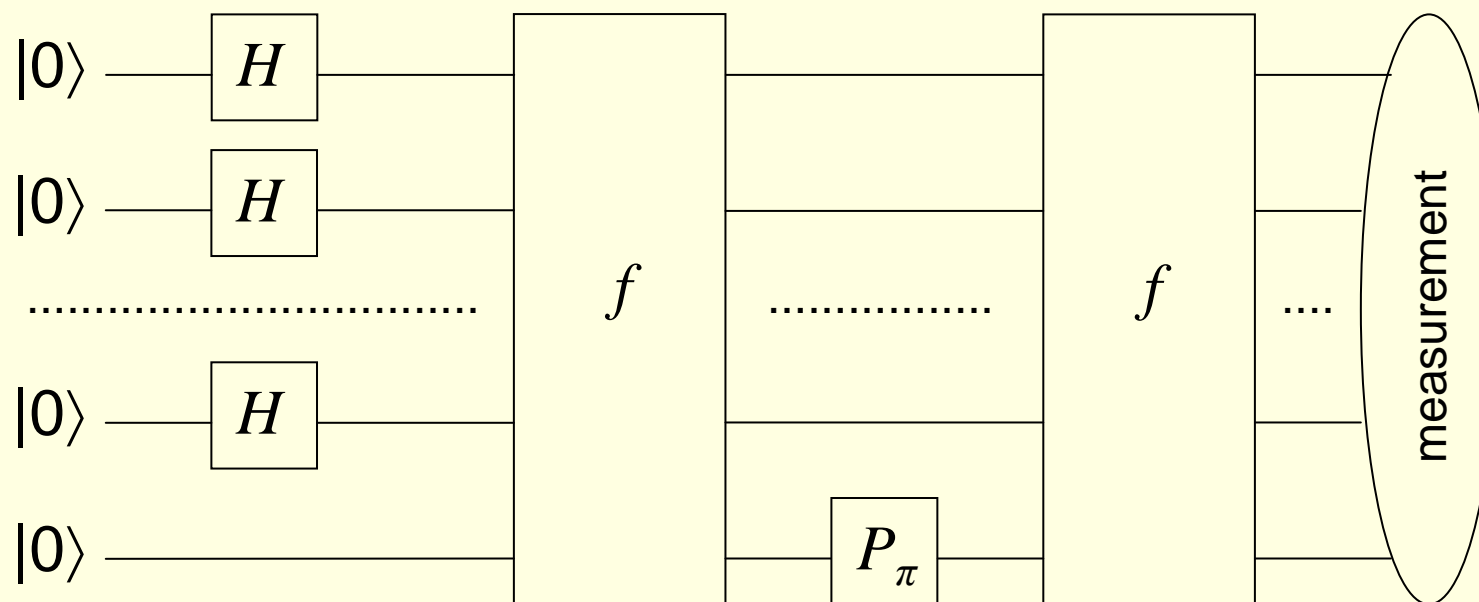
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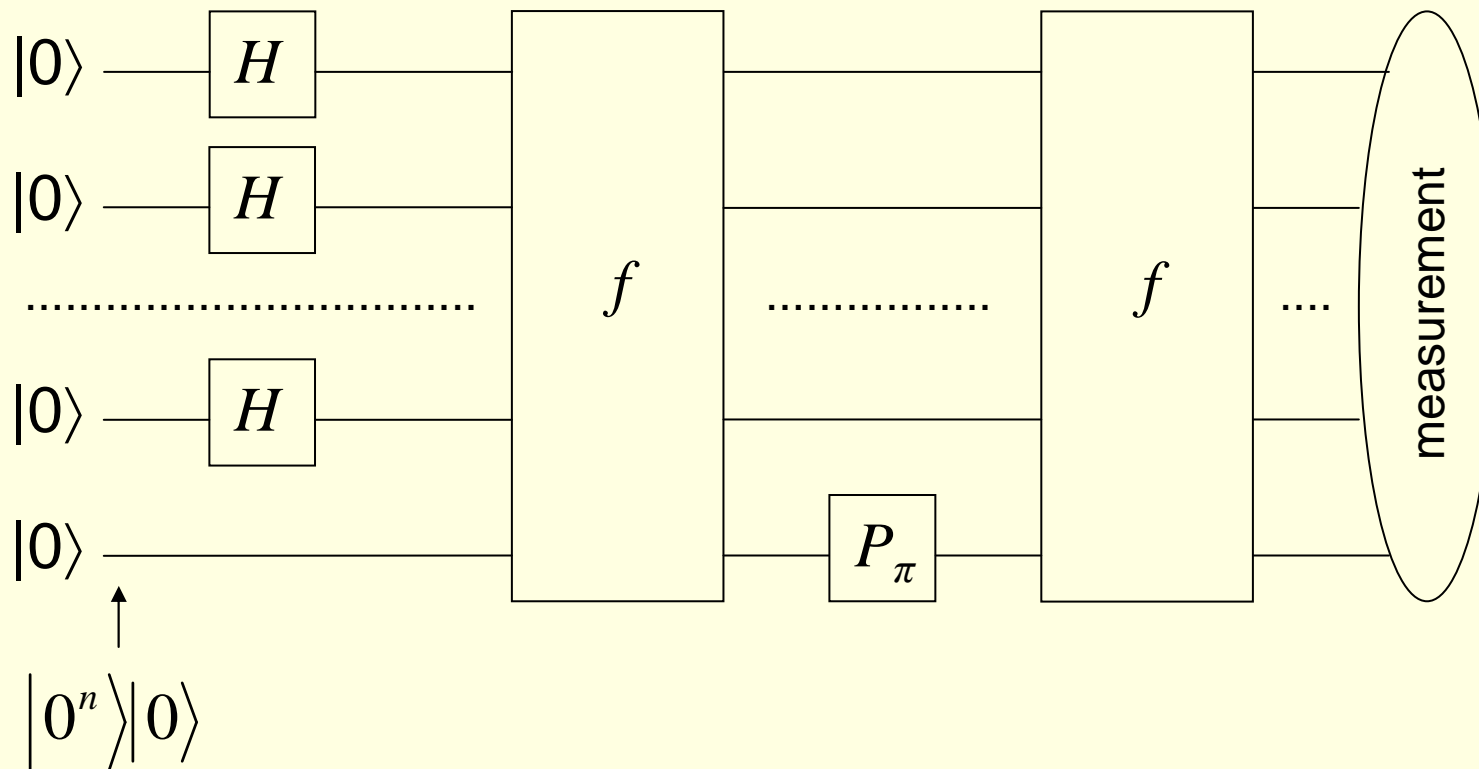
Efficient quantum algorithm

- Example: the Deutsch-Jozsa algorithm (1992)



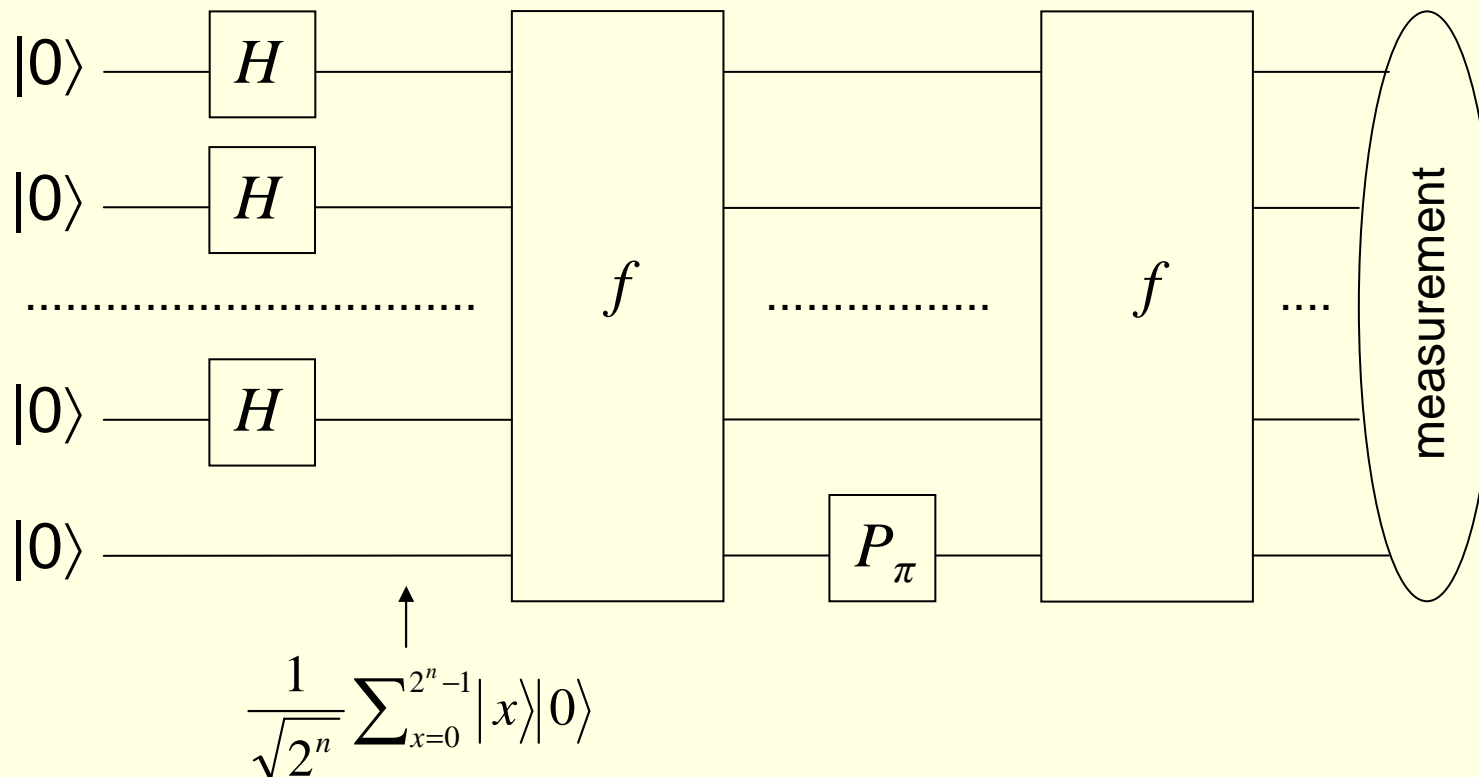
Efficient quantum algorithm

■ Example: the Deutsch-Jozsa algorithm (1992)



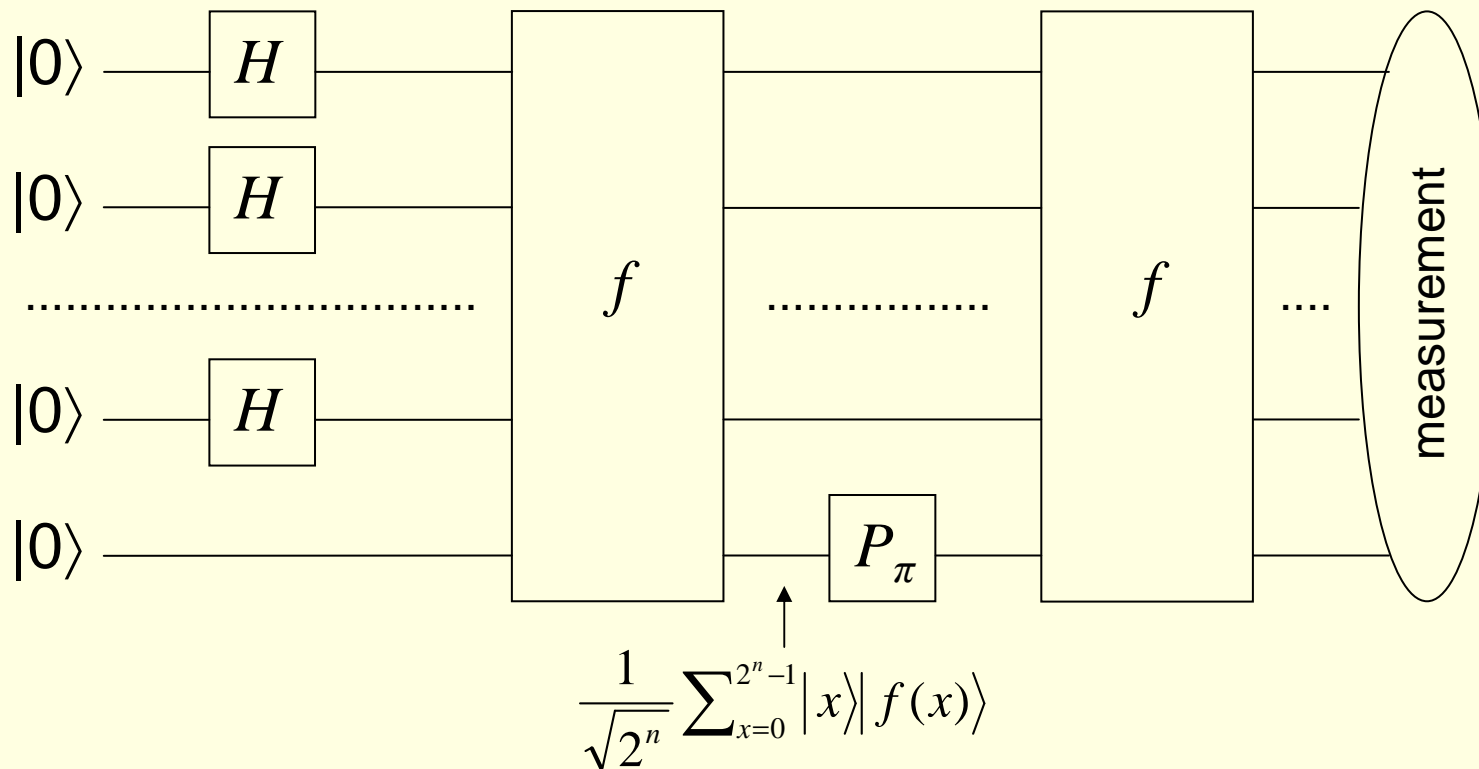
Efficient quantum algorithm

- Example: the Deutsch-Jozsa algorithm (1992)



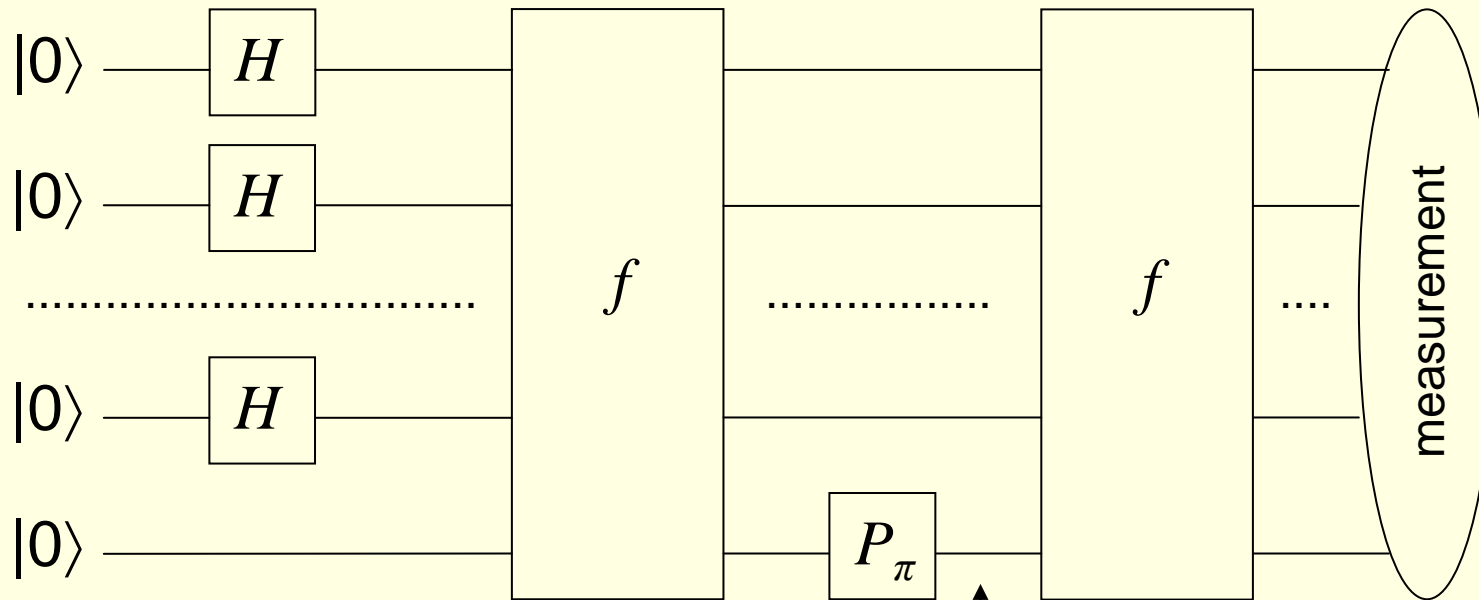
Efficient quantum algorithm

■ Example: the Deutsch-Jozsa algorithm (1992)



Efficient quantum algorithm

- Example: the Deutsch-Jozsa algorithm (1992)



$$\frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} (-1)^{f(x)} |x\rangle |f(x)\rangle$$



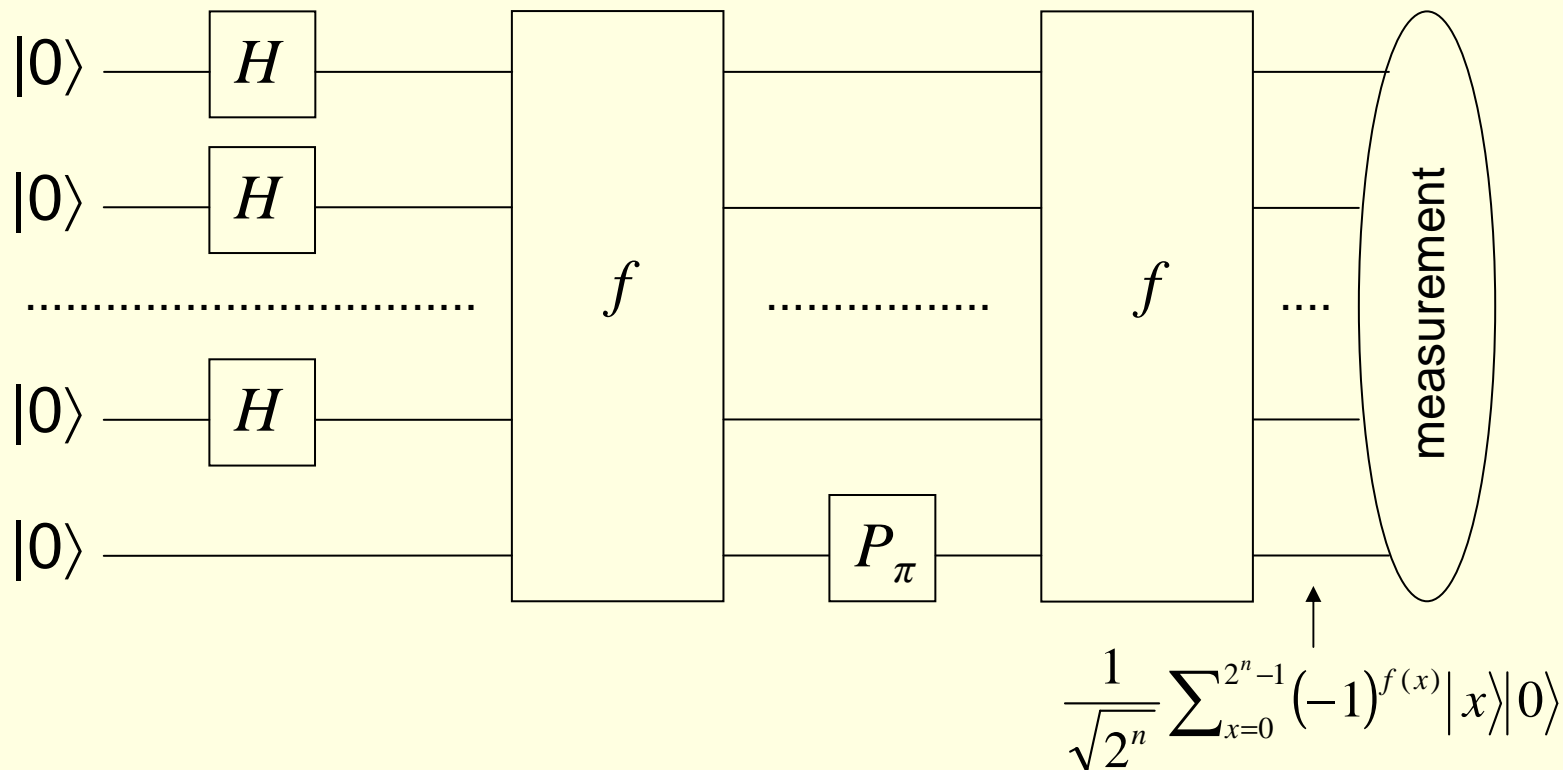
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Efficient quantum algorithm

- Example: the Deutsch-Jozsa algorithm (1992)



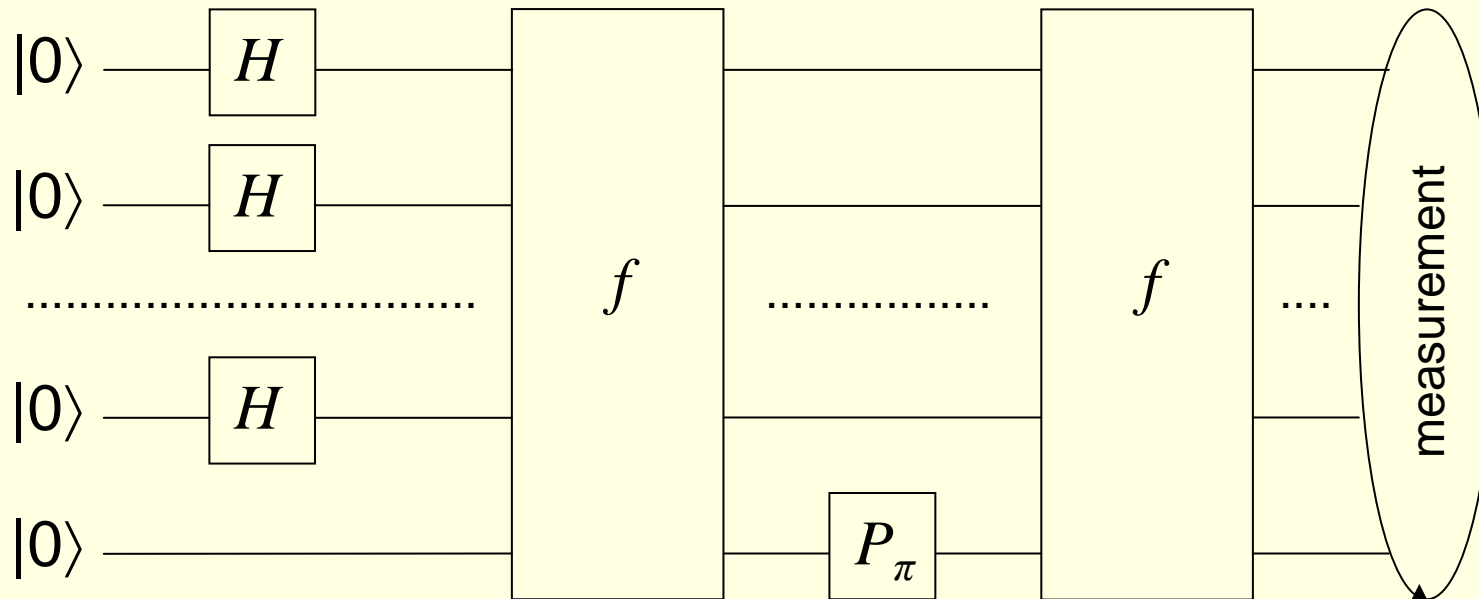
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Efficient quantum algorithm

■ Example: the Deutsch-Jozsa algorithm (1992)



$$|\psi\rangle = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} (-1)^{f(x)} |x\rangle |0\rangle, \quad E_{\text{constant}} = \text{span} \left(\frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle |0\rangle \right), \quad E_{\text{balanced}} = E_{\text{constant}}^\perp$$



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Efficient quantum algorithm

- E_{constant} is one-dimensional: $E_{\text{constant}} = \text{span}(|\theta\rangle)$ where

$$|\theta\rangle = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle |0\rangle$$

- Therefore the length of the projection of $|\psi\rangle$ on E_{constant} is the scalar product of $|\psi\rangle$ and $|\theta\rangle$, denoted by $\langle\psi|\theta\rangle$:

$$\langle\psi|\theta\rangle = \sum_{x=0}^{2^n-1} \frac{1}{\sqrt{2^n}} \cdot (-1)^{f(x)} \cdot \frac{1}{\sqrt{2^n}} = \frac{1}{2^n} \sum_{x=0}^{2^n-1} (-1)^{f(x)}$$

- ... which is by absolute value 1 iff f is constant, and 0 iff f is balanced
- The quantum algorithm is correct with probability 1 and with just two queries (classically $2^{n-1}+1$ are needed in the worst case)



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Efficient quantum algorithm

- Another example: Grover's algorithm (1996)
- Input: a black-box function $f: \{0,1\}^n \rightarrow \{0,1\}$ such that $f(x) = 1$ just for one (unknown) value, $x = x_0$; for all other values $f(x) = 0$
- The algorithm can query the black box; the number of queries determines the complexity of algorithm
- The black box queries on classical inputs work like this: input $|x\rangle$, output $(-1)^{f(x)}|x\rangle$
- Output: x_0



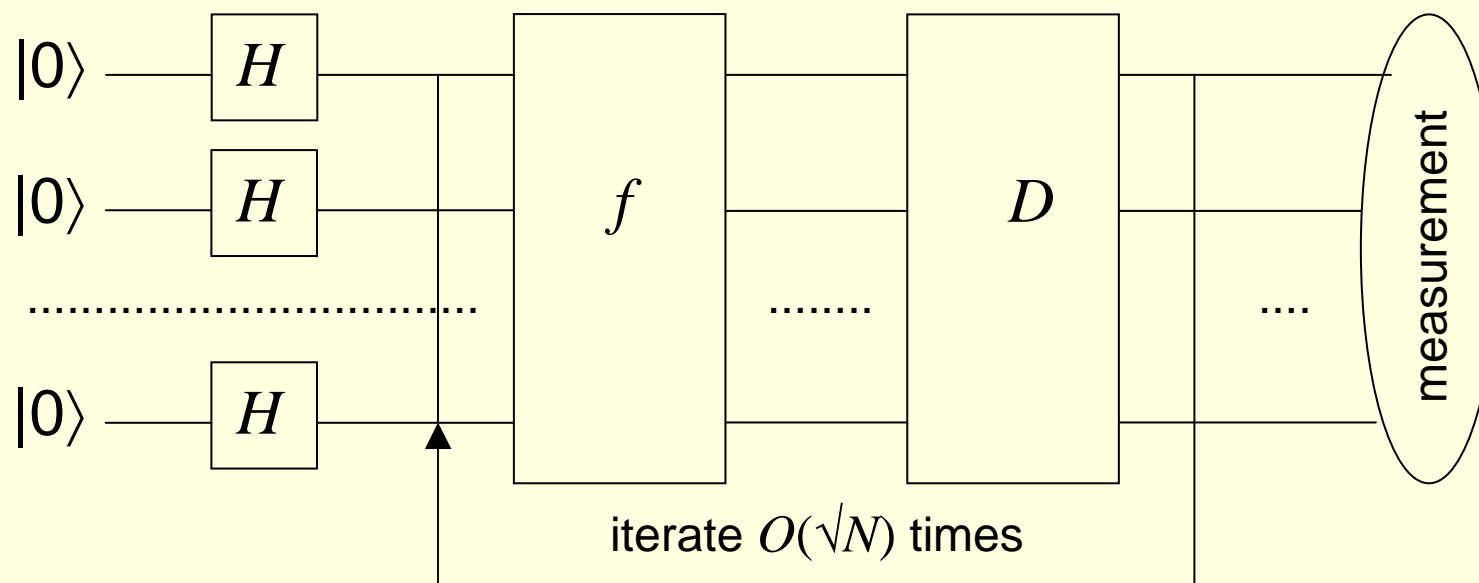
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Efficient quantum algorithm

- Example: the Grover's algorithm (1996)



$$N = 2^n$$

Efficient quantum algorithm

- The matrix of the f -query transformation

$$f\text{-query} = \begin{pmatrix} 1 & \dots & 0 & 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & 1 & 0 & 0 & \dots & 0 \\ 0 & \dots & 0 & -1 & 0 & \dots & 0 \\ 0 & \dots & 0 & 0 & 1 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & 0 & 0 & 0 & \dots & 1 \end{pmatrix} |x_0\rangle$$



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“Datorzinātnes pielietojumi un tās saiknes ar kvantu fiziku”

Nr.2009/0216/1DP/1.1.1.2.0/09/APIA/VIAA/044

Efficient quantum algorithm

- The matrix of the D (diffusion) transformation

$$D = \begin{pmatrix} \frac{2}{N} - 1 & \frac{2}{N} & \cdots & \frac{2}{N} \\ \frac{2}{N} & \frac{2}{N} - 1 & \cdots & \frac{2}{N} \\ \cdots & \cdots & \cdots & \cdots \\ \frac{2}{N} & \frac{2}{N} & \cdots & \frac{2}{N} - 1 \end{pmatrix} = \frac{2}{N} J - I$$

Efficient quantum algorithm

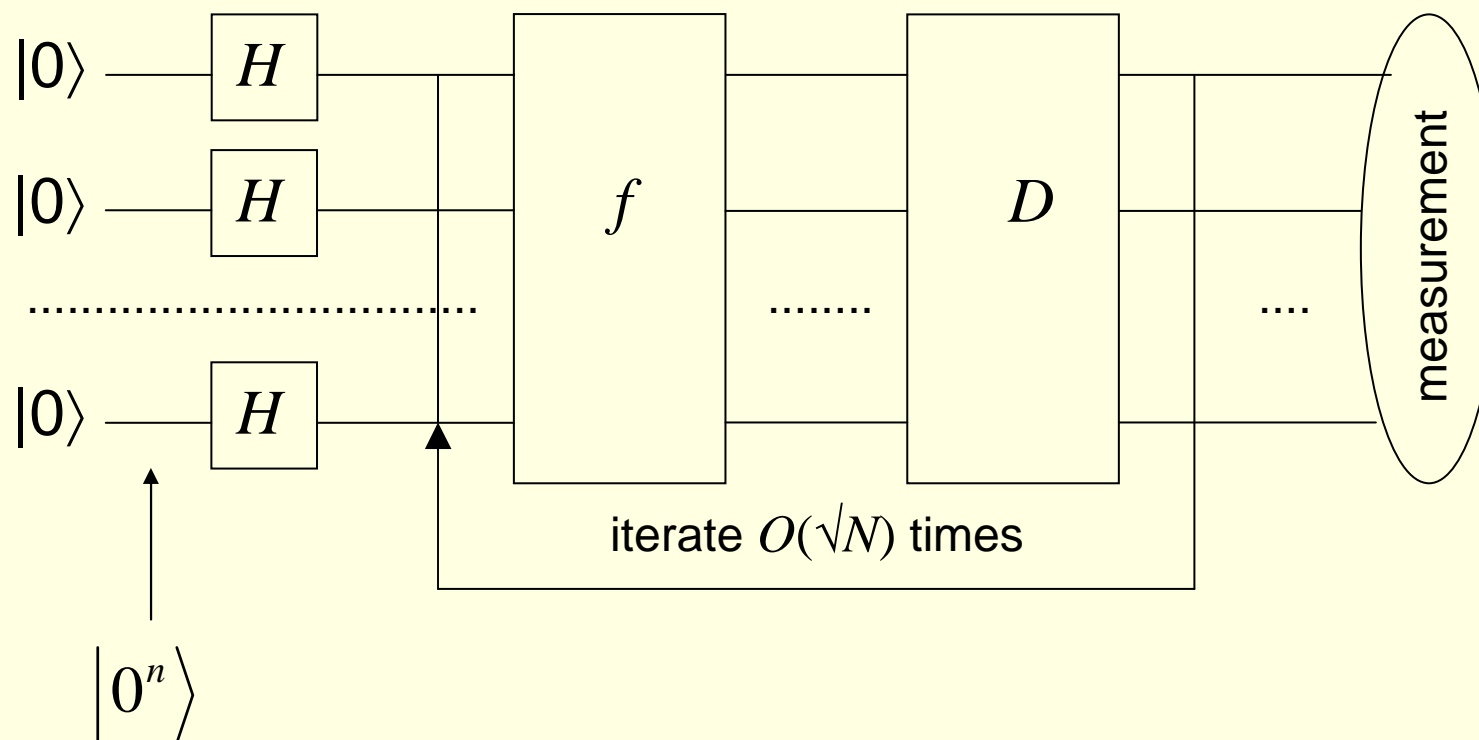
- D performs on the vector components inversion about their arithmetic mean

$$D \cdot \begin{pmatrix} a_1 \\ a_2 \\ \dots \\ a_N \end{pmatrix} = \left(\frac{2}{N} J - I \right) \cdot \begin{pmatrix} a_1 \\ a_2 \\ \dots \\ a_N \end{pmatrix} = \begin{pmatrix} 2m - a_1 \\ 2m - a_2 \\ \dots \\ 2m - a_N \end{pmatrix} = \begin{pmatrix} m + (m - a_1) \\ m + (m - a_2) \\ \dots \\ m + (m - a_N) \end{pmatrix}$$

where $m = \frac{a_1 + a_2 + \dots + a_N}{N}$

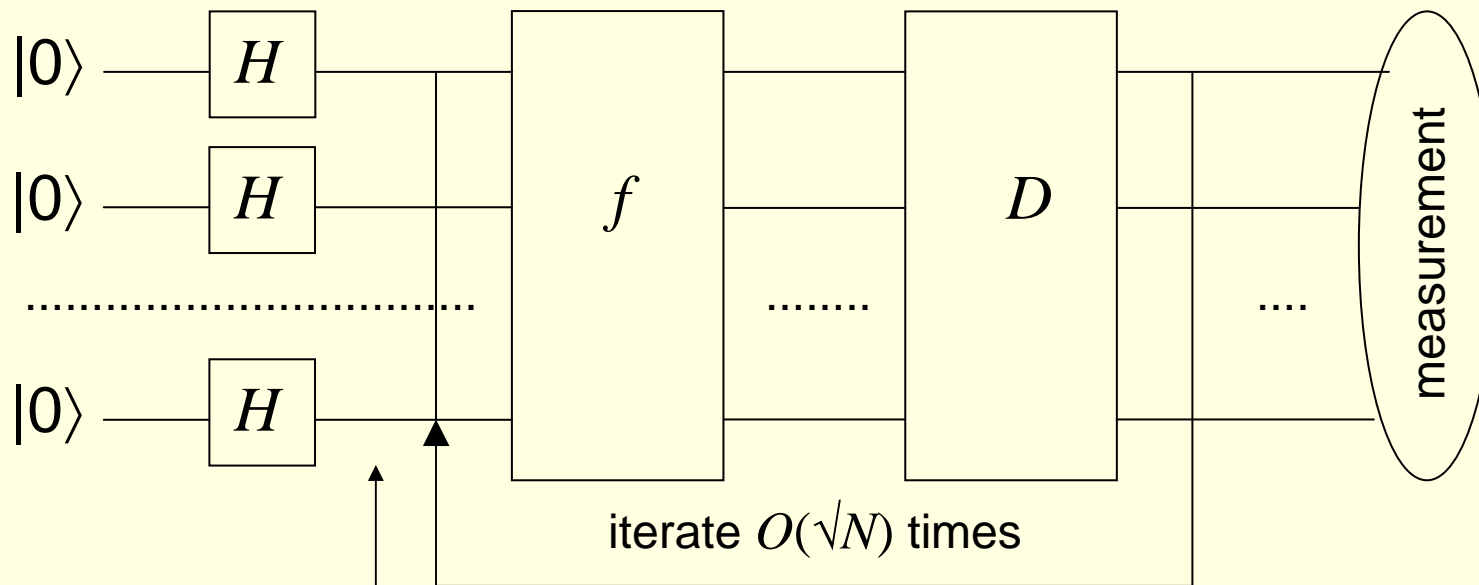
Efficient quantum algorithm

■ Example: the Grover's algorithm (1996)



Efficient quantum algorithm

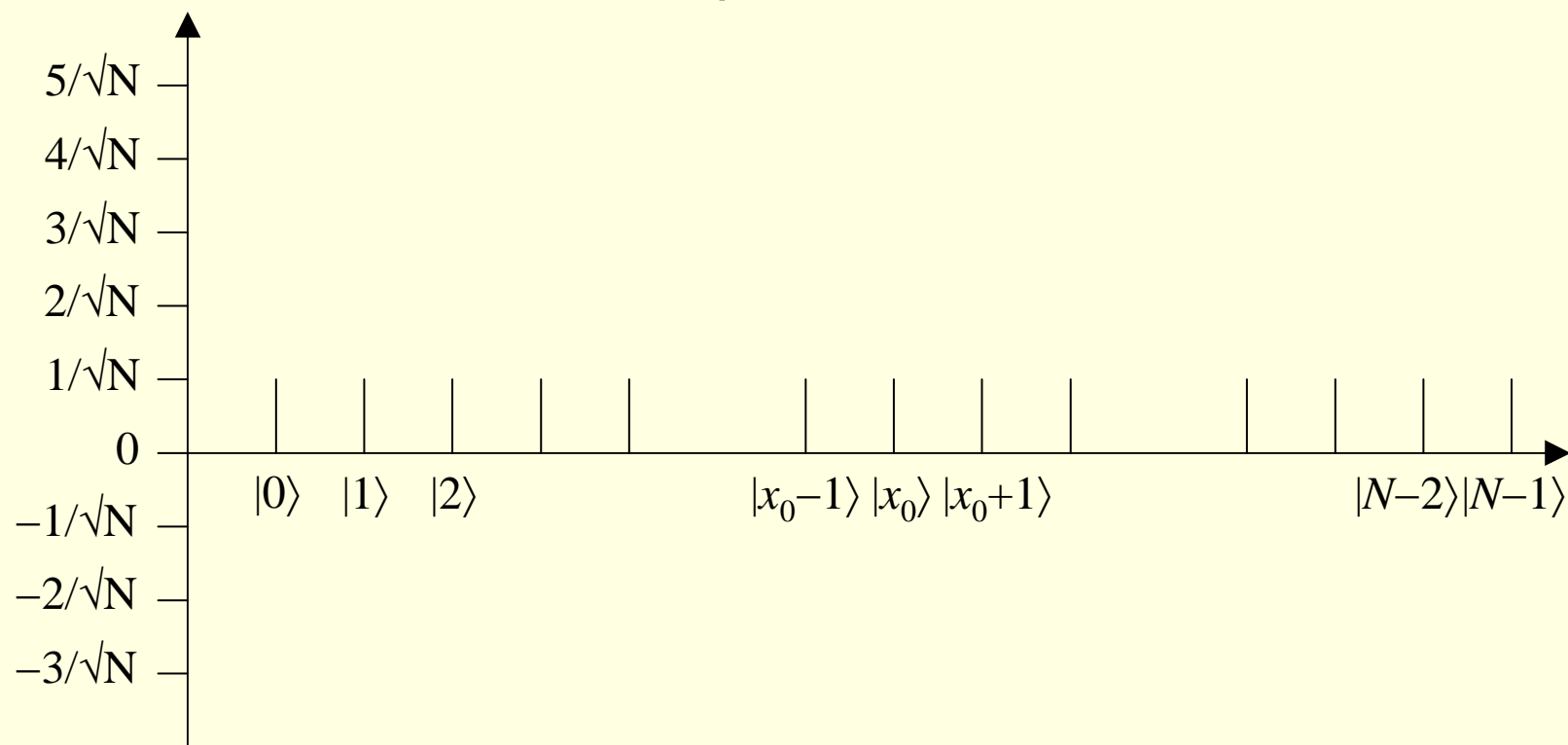
- Example: the Grover's algorithm (1996)



$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle = \left(\frac{1}{\sqrt{N}} \quad \frac{1}{\sqrt{N}} \quad \dots \quad \frac{1}{\sqrt{N}} \right)^T$$

Efficient quantum algorithm

- The values of the components of the state vector



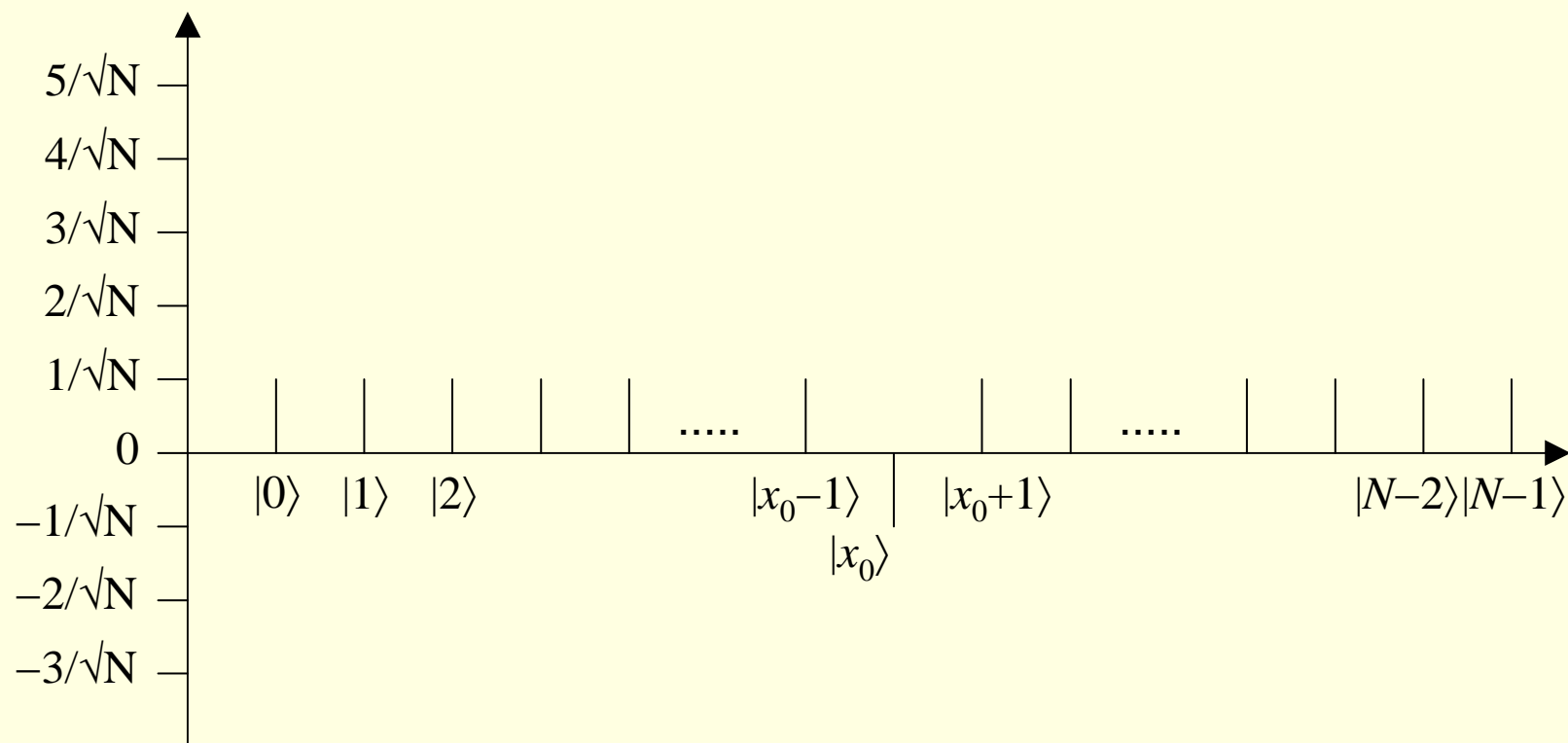
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Efficient quantum algorithm

■ After the f -query



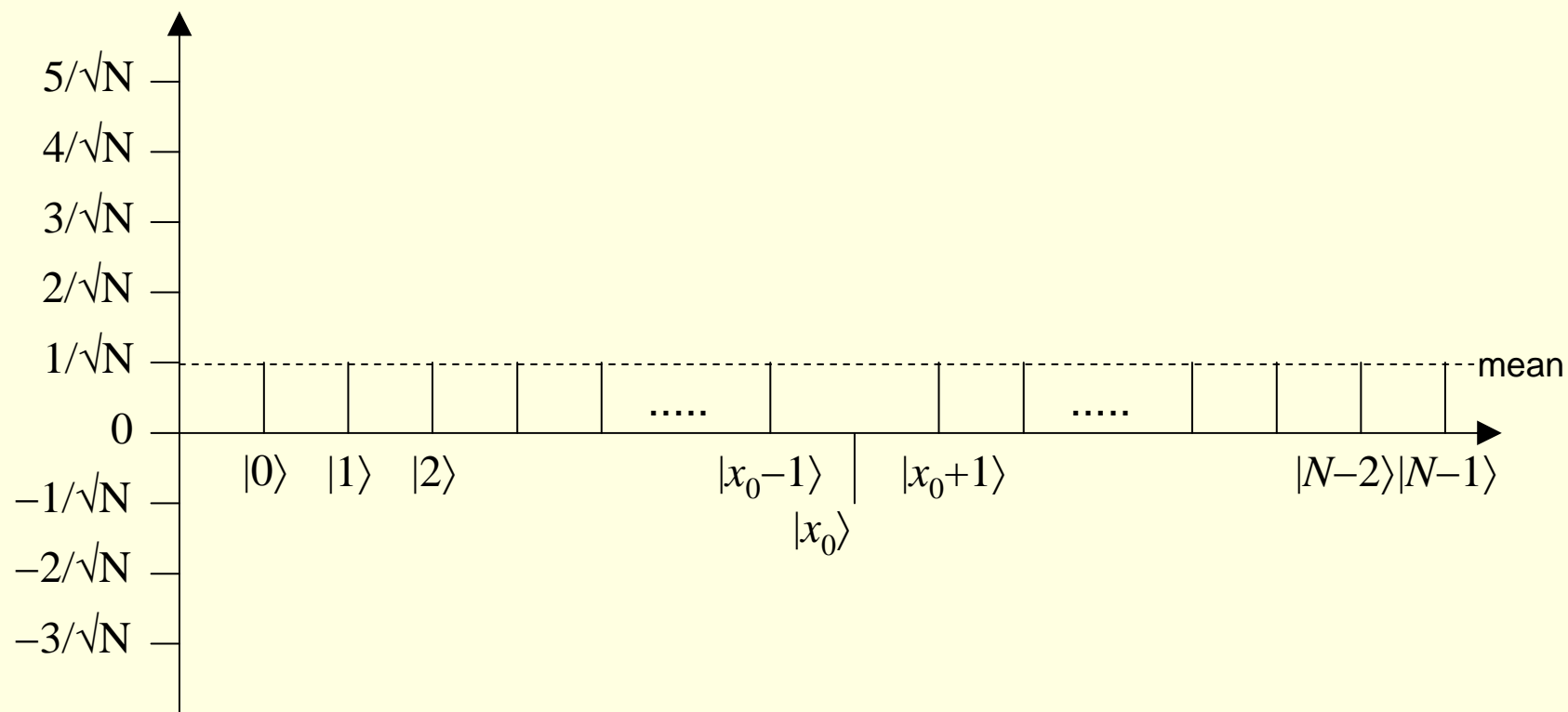
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Efficient quantum algorithm

■ After the f -query



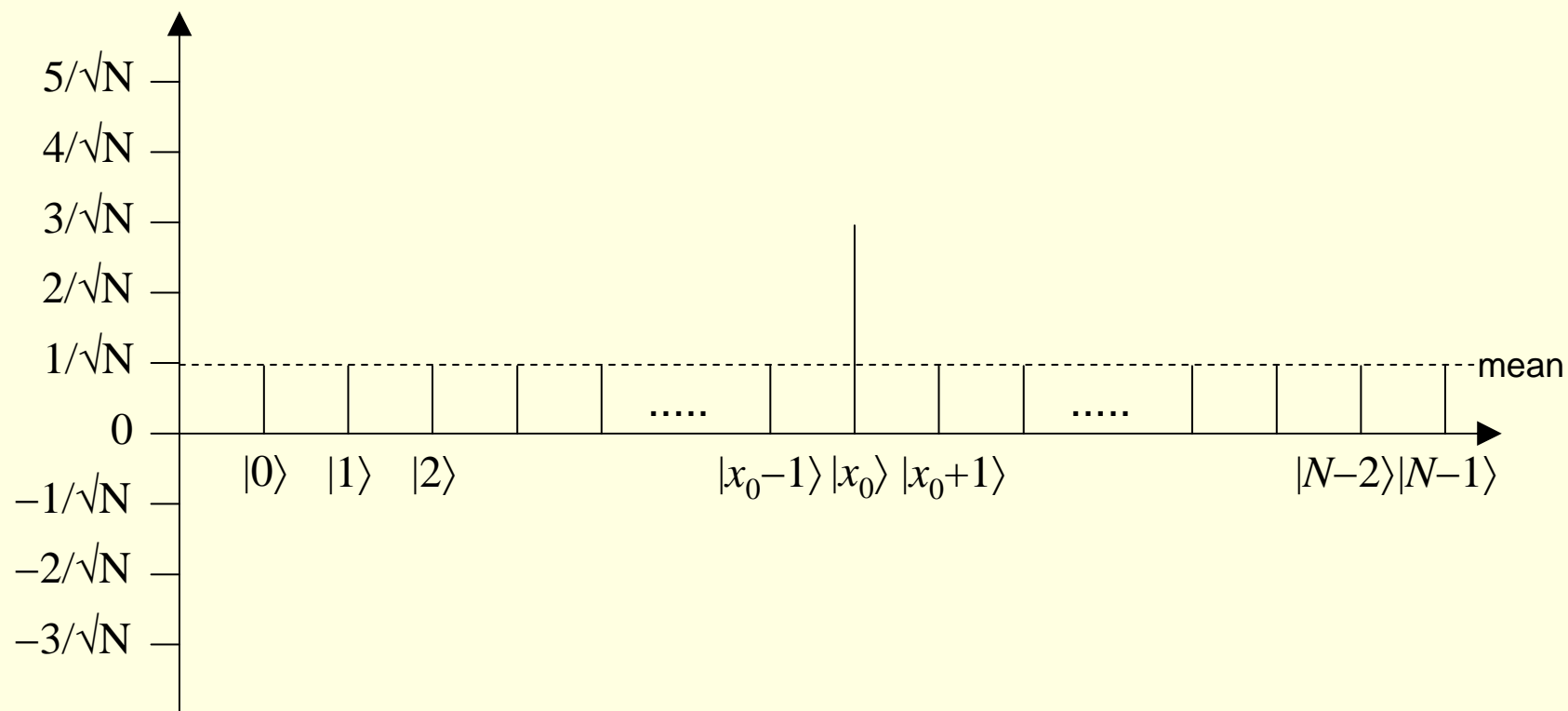
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Efficient quantum algorithm

■ After the diffusion



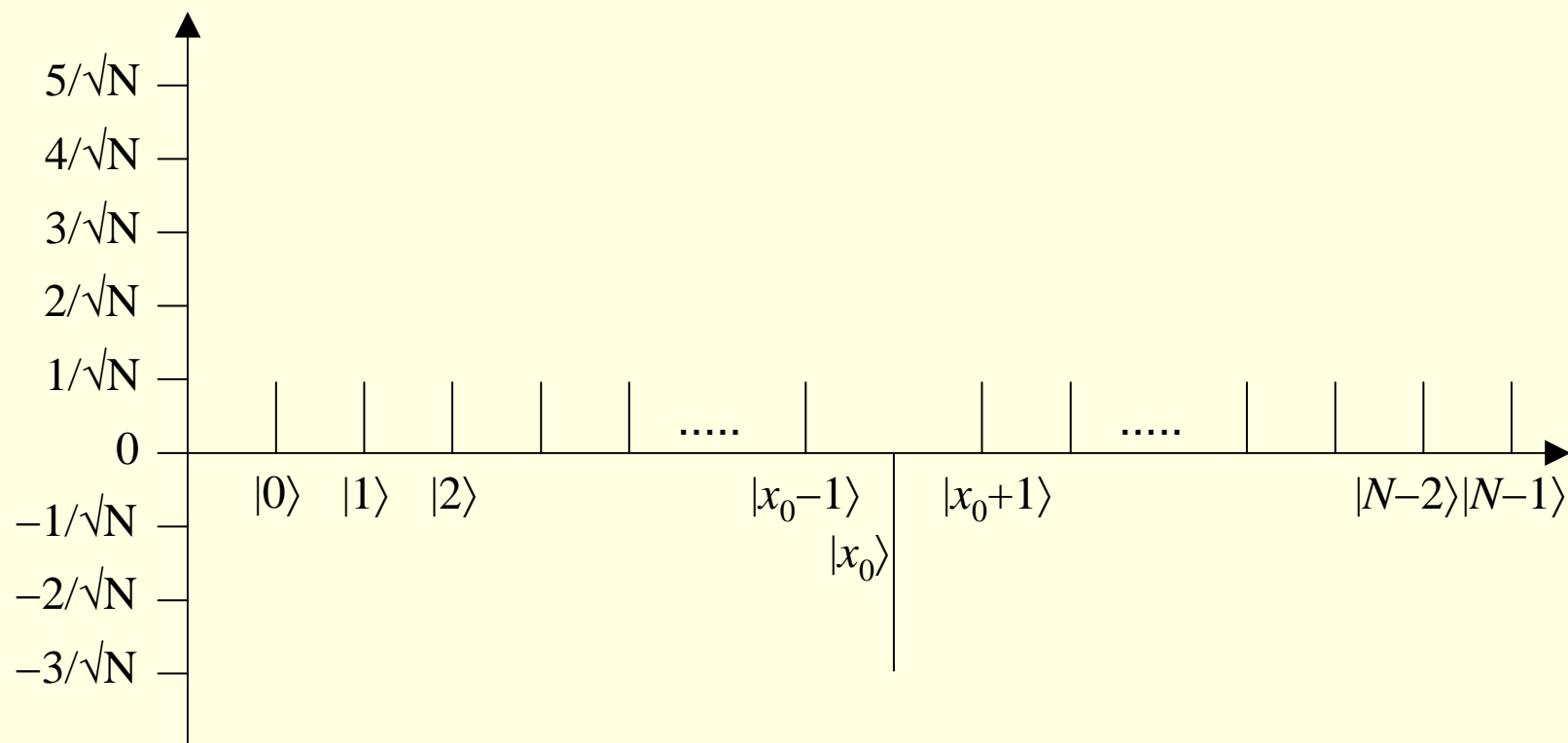
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Efficient quantum algorithm

■ After the f -query



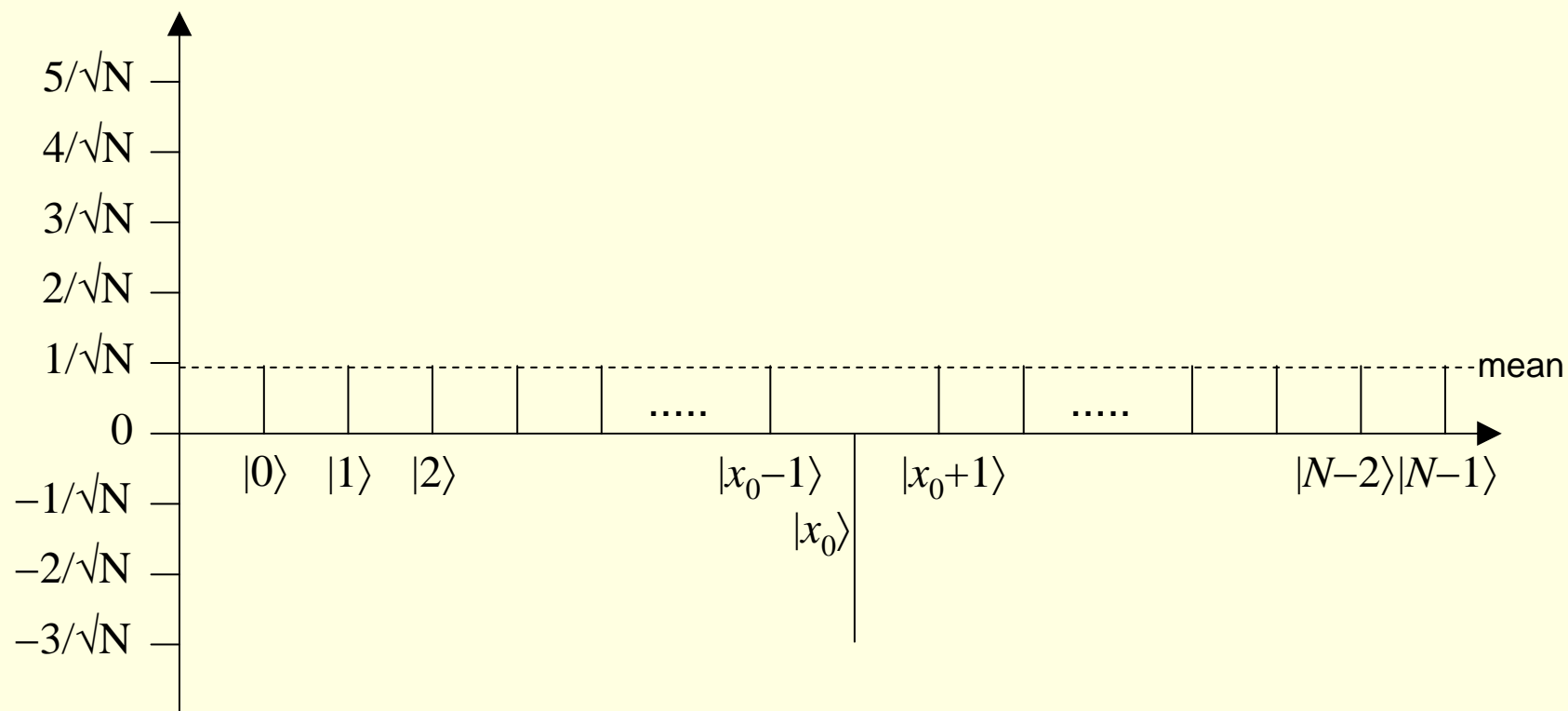
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Efficient quantum algorithm

■ After the f -query



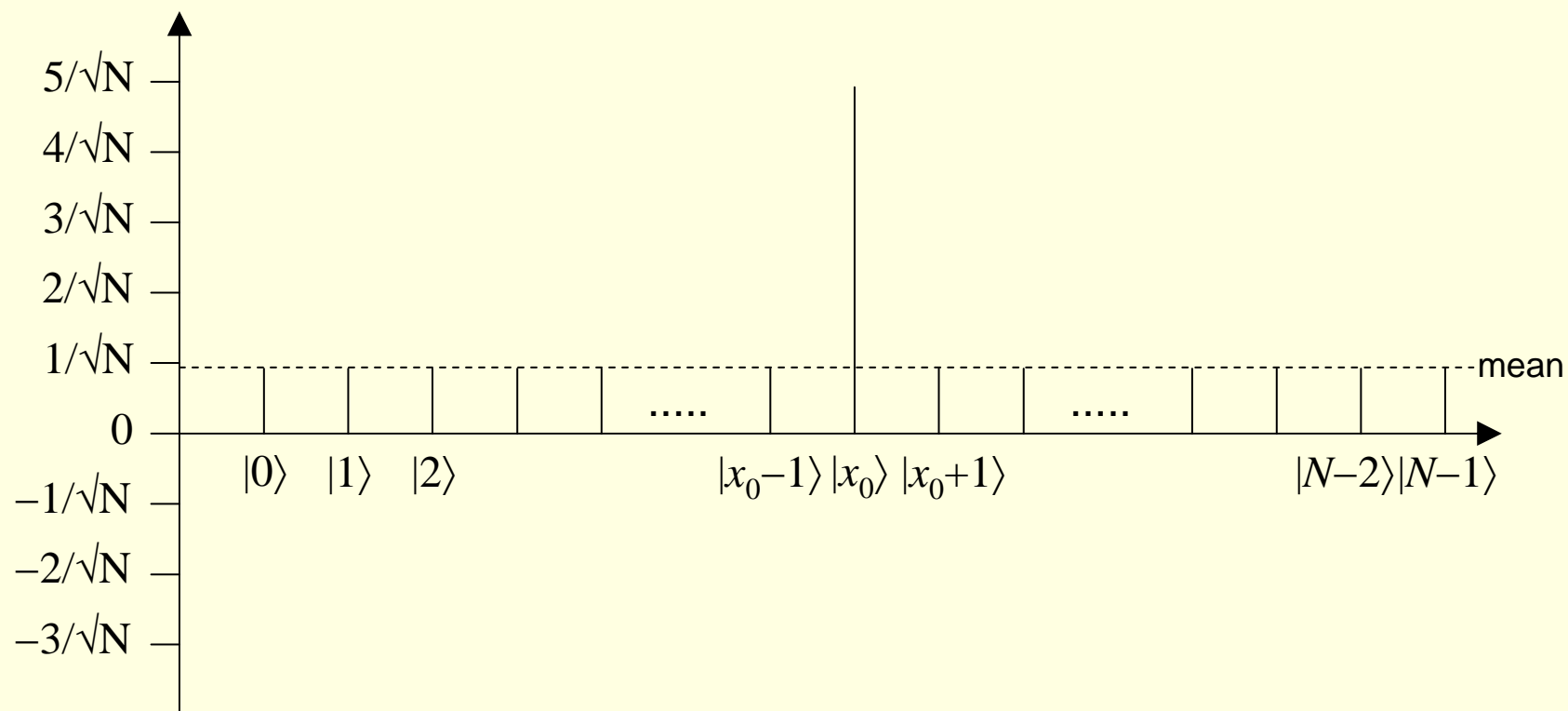
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Efficient quantum algorithm

■ After the diffusion



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“Datorzinātnes pielietojumi un tās saiknes ar kvantu fiziku”

Nr.2009/0216/1DP/1.1.1.2.0/09/APIA/VIAA/044

Efficient quantum algorithm

- After $O(\sqrt{N})$ iterations the amplitude at $|x_0\rangle$ practically reaches 1
- So the measurement at that moment gives $|x_0\rangle$ with probability (almost) 1
- Classically at least $O(N)$ queries are required



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“Datorzinātnes pielietojumi un tās saiknes ar kvantu fiziku”

Nr.2009/0216/1DP/1.1.1.2.0/09/APIA/VIAA/044

Thank you for your attention!

Questions?



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