

**“The Computational Complexity of Linear Optics”
by Scott Aaronson and Alex Arkhipov**

arXiv:1011.3245

Maris Ozols

April 21, 2011

The BIG dilemma...

The great dream

Quantum computers are more powerful than classical computers

The sad possibility

Church–Turing thesis: Everything that is efficiently computable by any physical device is efficiently computable by a Turing machine

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Polynomial
hierarchy
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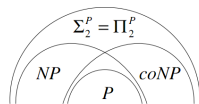
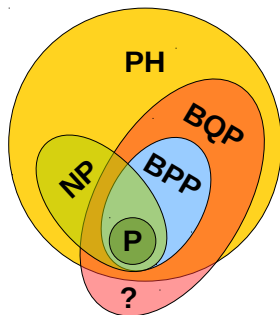


PHOTO: The scene of the devastation.

Complexity theory crash course



Complexity classes

- ▶ **P** – polynomial time
- ▶ **NP** – non-deterministic polynomial time
- ▶ **PH** – polynomial hierarchy (2nd order logic)
- ▶ **BPP** – bounded-error probabilistic polynomial time
- ▶ **BQP** – bounded-error quantum polynomial time

Computation with non-interacting bosons

Model of computation

- ▶ **Parameters:** n photons in $m = \text{poly}(n)$ modes
- ▶ **State space:** $\text{span}\{|s_1, \dots, s_m\rangle : s_k \geq 0, \sum_{k=1}^m s_k = n\}$
- ▶ **Initial state:** $|1_n\rangle := |1, \dots, 1, 0, \dots, 0\rangle$
- ▶ **Transformations:** $\varphi_n(U)$ for any $U \in \text{U}(m)$, where φ_n extends the action from 1 to n photons
- ▶ **Measurement:** the number of photons in each mode

The BOSONSAMPLING problem

Given a description of $U \in \text{U}(m)$, produce samples from the probability distribution

$$\Pr[S] := |\langle S | \varphi_n(U) | 1_n \rangle|^2$$

Transition matrix

Definition

If $|S\rangle = |s_1, \dots, s_m\rangle$ and $|T\rangle = |t_1, \dots, t_m\rangle$ then φ_n is defined as

$$\langle S|\varphi_n(A)|T\rangle = \frac{\text{perm}(A_{S,T})}{\sqrt{s_1! \cdots s_m! t_1! \cdots t_m!}}$$

where $A_{S,T}$ is the $n \times n$ matrix obtained by taking s_i copies of the i th row and t_j copies of j th column of A

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Example

$$\varphi_2 : \begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{matrix} |10\rangle \\ |01\rangle \end{matrix} \mapsto \begin{pmatrix} a^2 & \sqrt{2}ab & b^2 \\ \sqrt{2}ac & ad + bc & \sqrt{2}bd \\ c^2 & \sqrt{2}cd & d^2 \end{pmatrix} \begin{matrix} |20\rangle \\ |11\rangle \\ |02\rangle \end{matrix}$$

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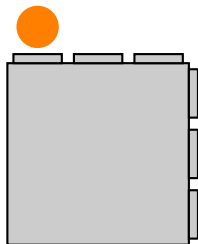
Properties

- ▶ φ_n is a homomorphism: $\varphi_n(A \cdot B) = \varphi_n(A) \cdot \varphi_n(B)$
- ▶ if U is unitary then so is $\varphi_n(U)$

Permanent

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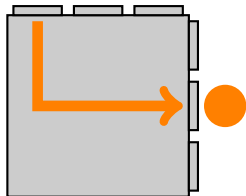
The magic box



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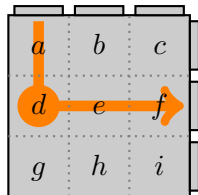
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<i>a</i>	<i>b</i>	<i>c</i>
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Permanent

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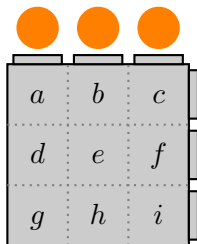


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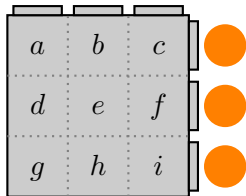
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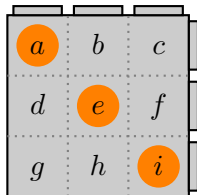
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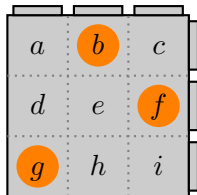
aei

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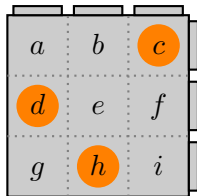


$$aei + gbf$$

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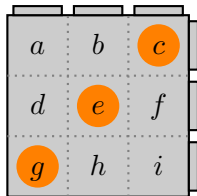


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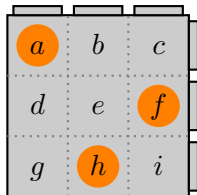


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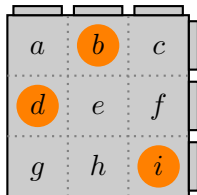


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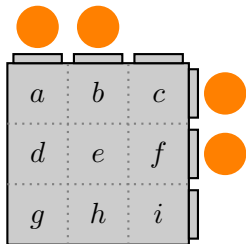
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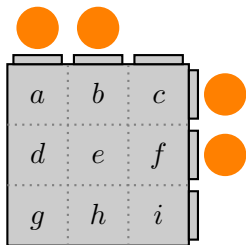
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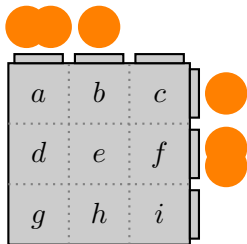
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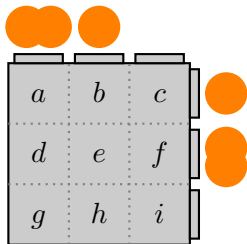
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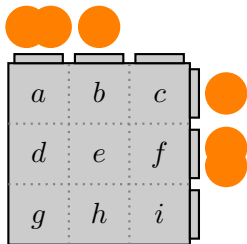
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Theorem (Valiant '79)

Computing $\text{perm}(A)$ is #P-hard

Main results

Theorem (exact case)

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Theorem (approximate case)

If the following two conjectures are true:

1. the permanent of a random Gaussian matrix is $\#P$ -hard to approximate and
2. it is not too concentrated around 0

then it is not possible to approximately solve the `BOSONSAMPLING` problem, unless polynomial hierarchy collapses

Experimental feasibility

Linear optics

- ▶ prepare photons using single photon sources
- ▶ use beam splitters and phase shifters to implement U
- ▶ use photodetectors to perform the readout
- ▶ the order of parameters: $n = 10$, $m = 20$

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Problems

- ▶ need good photon sources and detectors
- ▶ all n photons must arrive at the destination at the same time

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